

# Pre-Defined Functions

- $\text{bot} :: a$
- $\text{isa}_{\underline{\text{constr}}} :: \underline{\text{type}} \rightarrow \text{Bool}$   
for all  $\underline{\text{constr}} \in \text{Con}$ , where  $\underline{\text{constr}} :: \underline{\text{type}}_1 \rightarrow \dots \rightarrow \underline{\text{type}}_n \rightarrow \underline{\text{type}}$
- $\text{argof}_{\underline{\text{constr}}} :: \underline{\text{type}} \rightarrow (\underline{\text{type}}_1, \dots, \underline{\text{type}}_n)$   
for all  $\underline{\text{constr}} \in \text{Con}$ , wobei  $\underline{\text{constr}} :: \underline{\text{type}}_1 \rightarrow \dots \rightarrow \underline{\text{type}}_n \rightarrow \underline{\text{type}}$
- $\text{isa}_{n\text{-tuple}} :: (a_1, \dots, a_n) \rightarrow \text{Bool}$   
for all  $n \in \{0, 2, 3, \dots, k\}$
- $\text{sel}_{n,i} :: (a_1, \dots, a_n) \rightarrow a_i$   
for all  $2 \leq n \leq k, 1 \leq i \leq n$ .

```
append Nil z = z
```

```
append (Cons x y) z = Cons x (append y z)
```

↓ (1)

```
append = \x1 x2 -> case (x1, x2) of (Nil, z) -> z  
                                (Cons x y, z) -> Cons x (append y z)
```

↓ (2)

```
append = \x1 ->  
        (\x2 -> case (x1, x2) of (Nil, z) -> z  
                                (Cons x y, z) -> Cons x (append y z)).
```

↓ (4)

```
append = \x1 ->  
        (\x2 -> match (Nil, z) (x1, x2) z  
                            (match (Cons x y, z) (x1, x2) (Cons x (append y z)) bot))
```

```

append = \x1 ->
  (\x2 -> match (Nil, z) (x1, x2) z
    (match (Cons x y, z) (x1, x2) (Cons x (append y z)) bot))
    ↓ (9)
append = \x1 ->
  (\x2 ->
    match Nil
      (sel2,1 (x1, x2))
      (match z
        (sel2,2 (x1, x2))
        z
        (match (Cons x y)
          (sel2,1 (x1, x2))
          (match z
            (sel2,2 (x1, x2))
            (Cons x (append y z))
            bot)
            bot))
        (match (Cons x y)
          (sel2,1 (x1, x2))
          (match z
            (sel2,2 (x1, x2))
            (Cons x (append y z))
            bot)
            bot))
        bot))

```

```
append = \x1 -> (\x2 -> match Nil
                      x1
                      (match z
                        x2
                        z
                        (match (Cons x y)
                          :
                          )
                        (match (Cons x y)
                          x1
                          (match z
                            x2
                            (Cons x (append y z))
                            bot)
                          bot)))

```

↓ (5)

```
append = \x1 -> (\x2 -> match Nil
                      x1
                      (\z -> z) x2
                      (match (Cons x y)
                        x1
                        (\z -> (Cons x (append y z))) x2
                        bot))
```

```
append = \x1 ->
  (\x2 ->
    match Nil x1 x2
      (match (Cons x y) x1 (Cons x (append y x2)) bot))
```

↓ (7)

```
append = \x1 ->
  (\x2 ->
    if (isaNil x1)
      then (match () (argofNil x1) x2
        (if (isaCons x1)
          then (match (x, y)
            (argofCons x1)
            (Cons x (append y x2))
            bot)
          else bot))
      else (if (isaCons x1)
        then (match (x, y)
          (argofCons x1)
          (Cons x (append y x2))
          bot)
        else bot))
```

↓ (8)

```

append = \x1 -> (\x2 ->
  if (isaNil x1)
  then (if (isa() (argofNil x1))
         then x2
         else (if (isaCons x1)
                then (match (x, y)
                       (argofCons x1)
                       (Cons x (append y x2))
                       bot)
                else bot))
  else (if (isaCons x1)
         then (match (x, y) (argofCons x1) (Cons x (append y x2)) bot)
         else bot))

```

↓ (9), (5)

```

append = \x1 -> (\x2 ->
  if (isaNil x1)
  then x2
  else (if (isaCons x1)
         then (Cons (sel2,1 (argofCons x1))
                     (append (sel2,2 (argofCons x1)) x2)))
         else bot))

```

# Transformation Into Simple HASKELL-Programs

## (1) Transformation of function declarations into pattern declarations

$$\frac{\underline{\text{var}} \underline{\text{pat}}_1^1 \dots \underline{\text{pat}}_n^1 = \underline{\text{exp}}^1; \dots; \underline{\text{var}} \underline{\text{pat}}_1^k \dots \underline{\text{pat}}_n^k = \underline{\text{exp}}^k}{\underline{\text{var}} = \lambda x_1 \dots x_n \rightarrow \text{case } (x_1, \dots, x_n) \text{ of } \{(\underline{\text{pat}}_1^1, \dots, \underline{\text{pat}}_n^1) \rightarrow \underline{\text{exp}}^1; \\ \vdots \\ (\underline{\text{pat}}_1^k, \dots, \underline{\text{pat}}_n^k) \rightarrow \underline{\text{exp}}^k\}}$$

where  $x_1, \dots, x_n$  must be new variables,  $n > 0$ ,  
and these are all declarations for var

## (2) Transforming lambda-expressions with several patterns

$$\frac{\lambda \underline{\text{pat}}_1 \dots \underline{\text{pat}}_n \rightarrow \underline{\text{exp}}}{\underline{\text{pat}}_1 \rightarrow (\lambda \underline{\text{pat}}_2 \rightarrow \dots (\lambda \underline{\text{pat}}_n \rightarrow \underline{\text{exp}}) \dots)} \quad \text{where } n \geq 2$$

### (3) Transforming lambda-patterns into case

$$\frac{\backslash \underline{\text{pat}} \rightarrow \underline{\text{exp}}}{\backslash \underline{\text{var}} \rightarrow \text{case } \underline{\text{var}} \text{ of } \underline{\text{pat}} \rightarrow \underline{\text{exp}}}$$

where pat is no variable, var is new

### (4) Transforming case into match

$$\frac{\text{case } \underline{\text{exp}} \text{ of } \{ \underline{\text{pat}}_1 \rightarrow \underline{\text{exp}}_1; \dots; \underline{\text{pat}}_n \rightarrow \underline{\text{exp}}_n \}}{\begin{aligned} &\text{match } \underline{\text{pat}}_1 \underline{\text{exp}} \underline{\text{exp}}_1 \\ &\quad (\text{match } \underline{\text{pat}}_2 \underline{\text{exp}} \underline{\text{exp}}_2 \\ &\quad \dots \\ &\quad (\text{match } \underline{\text{pat}}_n \underline{\text{exp}} \underline{\text{exp}}_n \text{ bot}) \dots) \end{aligned}}$$

### (5) match of variables

$$\frac{\text{match } \underline{\text{var}} \underline{\text{exp}} \underline{\text{exp}}_1 \underline{\text{exp}}_2}{(\backslash \underline{\text{var}} \rightarrow \underline{\text{exp}}_1) \underline{\text{exp}}}$$

## (6) match of joker pattern

$$\frac{\text{match } \underline{\text{exp}} \text{ } \underline{\text{exp}_1} \text{ } \underline{\text{exp}_2}}{\underline{\text{exp}_1}}$$

## (7) match of constructors

$$\frac{\text{match } (\underline{\text{constr}} \text{ } \underline{\text{pat}_1} \dots \underline{\text{pat}_n}) \text{ } \underline{\text{exp}} \text{ } \underline{\text{exp}_1} \text{ } \underline{\text{exp}_2}}{\text{if } (\text{isa}_{\underline{\text{constr}}} \text{ } \underline{\text{exp}}) \text{ then } (\text{match } (\underline{\text{pat}_1}, \dots, \underline{\text{pat}_n}) \\ (\text{argof}_{\underline{\text{constr}}} \text{ } \underline{\text{exp}}) \text{ } \underline{\text{exp}_1} \text{ } \underline{\text{exp}_2}) \text{ else } \underline{\text{exp}_2}}$$

## (8) match of empty tuple

$$\frac{\text{match } () \text{ } \underline{\text{exp}} \text{ } \underline{\text{exp}_1} \text{ } \underline{\text{exp}_2}}{\text{if } (\text{isa}_{0\text{-tuple}} \text{ } \underline{\text{exp}}) \text{ then } \underline{\text{exp}_1} \text{ else } \underline{\text{exp}_2}}$$

## (9) match of non-empty tuple

$$\frac{\text{match } (\underline{\text{pat}}_1, \dots, \underline{\text{pat}}_n) \text{ exp } \underline{\text{exp}}_1 \underline{\text{exp}}_2}{\begin{aligned} &\text{if } (\text{isa}_{n-\text{tuple}} \underline{\text{exp}}) \\ &\text{then match } \underline{\text{pat}}_1 (\underline{\text{sel}}_{n,1} \underline{\text{exp}}) \\ &\quad (\text{match } \underline{\text{pat}}_2 (\underline{\text{sel}}_{n,2} \underline{\text{exp}}) \\ &\quad \dots (\text{match } \underline{\text{pat}}_n (\underline{\text{sel}}_{n,n} \underline{\text{exp}}) \underline{\text{exp}}_1 \underline{\text{exp}}_2) \dots \\ &\quad \underline{\text{exp}}_2) \underline{\text{exp}}_2 \\ &\text{else } \underline{\text{exp}}_2 \end{aligned}} \quad \text{where } n \geq 2$$

## (10) Separation of Declarations

$$\frac{\text{let } P \text{ in } \underline{\text{exp}}}{\begin{aligned} &\text{let } P_1 \text{ in } \\ &\quad \text{let } P_2 \text{ in } \\ &\quad \dots \\ &\quad \text{let } P_k \text{ in } \underline{\text{exp}} \end{aligned}} \quad \text{where } P_1, \dots, P_k \text{ is a separation} \\ \text{of the declaration } P$$

## (11) Transforming sequences of declarations into a single declaration

$$\frac{\{\underline{\text{var}}_1 = \underline{\text{exp}}_1; \dots; \underline{\text{var}}_n = \underline{\text{exp}}_n\}}{(\underline{\text{var}}_1, \dots, \underline{\text{var}}_n) = (\underline{\text{exp}}_1, \dots, \underline{\text{exp}}_n)}$$

where  $n \geq 2$ ,  $\underline{\text{var}}_i \neq \underline{\text{var}}_j$  for  $i \neq j$  and  $\underline{\text{var}}_1 \sim_P \dots \sim_P \underline{\text{var}}_n$ . Here,  $P = \{\underline{\text{var}}_1 = \underline{\exp}_1; \dots; \underline{\text{var}}_n = \underline{\exp}_n\}$ .

## (12) Declaration of several variables

$$\frac{\text{let } (\underline{\text{var}}_1, \dots, \underline{\text{var}}_n) = \underline{\text{exp}} \text{ in } \underline{\text{exp}}'}{\text{let } \underline{\text{var}} = \text{match } (\underline{\text{var}}_1, \dots, \underline{\text{var}}_n) \text{ (sel}_{n,1} \underline{\text{var}}, \dots, \text{sel}_{n,n} \underline{\text{var}}) \text{ exp' bot} \\ \text{in } \text{match } (\underline{\text{var}}_1, \dots, \underline{\text{var}}_n) \text{ var exp' bot}}$$

where  $n \geq 2$  and var is a new variable