

Semantics by lub's of Approximations

`fact :: Int -> Int`

`fact = \x -> if x <= 0 then 1 else fact(x-1) * x`

Regard non-recursively defined approximations

`fact0 = \x -> bot`

`fact1 = \x -> if x <= 0 then 1 else fact0(x - 1) * x`

`fact2 = \x -> if x <= 0 then 1 else fact1(x - 1) * x`

...

Thus: `facti+1 = ff facti` where

`ff :: (Int -> Int) -> (Int -> Int)`

`ff g = \x -> if x <= 0 then 1 else g(x-1) * x`

Define semantics *fact* of `fact` as

$$fact = \sqcup \{ fact_0, fact_1, \dots \} = \sqcup \{ ff^i(\perp) \mid i \in \mathbb{N} \}$$

Semantics by Least Fixpoints

`fact :: Int -> Int`

`fact = \x -> if x <= 0 then 1 else fact(x-1) * x`

Semantics *fact* of `fact` should *satisfy* the defining equation

$$fact = \underbrace{\lambda x \rightarrow \text{if } x \leq 0 \text{ then } 1 \text{ else } fact(x-1) * x}_{ff(fact)}$$

where

`ff :: (Int -> Int) -> (Int -> Int)`

`ff g = \x -> if x <= 0 then 1 else g(x-1) * x`

fact should be undefined unless program enforces its definedness.

Define semantics *fact* of `fact` as

$$fact = \text{lfp } ff$$

Fixpoint Theorem

Both semantics are equivalent:

$$\text{lfp } ff = \sqcup \{ ff^i(\perp) \mid i \in \mathbb{N} \}$$

Thm. 2.1.17 (Fixpoint Theorem)

Let \sqsubseteq be a cpo on D and let $f : D \rightarrow D$ be continuous.

Then f has a least fixpoint and we have $\text{lfp } f = \sqcup \{ f^i(\perp) \mid i \in \mathbb{N} \}$.

Proof.

$\perp \sqsubseteq f(\perp) \sqsubseteq f^2(\perp) \sqsubseteq f^3(\perp) \sqsubseteq \dots$ by monotonicity of f

So $\{ f^i(\perp) \mid i \in \mathbb{N} \}$ is a chain

and $\sqcup \{ f^i(\perp) \mid i \in \mathbb{N} \}$ exists by completeness of \sqsubseteq .

To show: $\sqcup \{ f^i(\perp) \mid i \in \mathbb{N} \}$ is the least fixpoint of f