

# Better termination proving through cooperation

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# Termination Analysis: Invariants and Rank Functions

## Example

```
y := 1;  
while x > 0 do  
    x := x - y;  
    y := y + 1;  
done
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- Invariant  $y > 0$  and rank function  $x$  prove termination
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- How do we know that we need  $y > 0$ ?  $\curvearrowright$   $x$  requires it
- How do we know that  $x$  is a RF?  $\curvearrowright$   $y > 0$  proves it

# Termination by iterative strengthening: Idea

- 1 Safety: Provide samples (Counterexamples)
- 2 Rank tool: Find specific termination argument
- 3 Safety: Prove generality, or 1

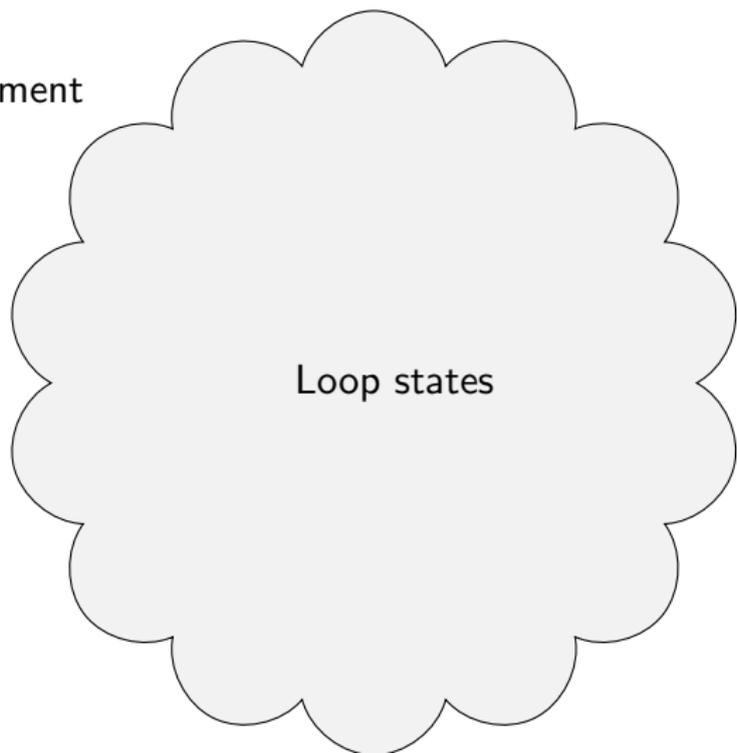
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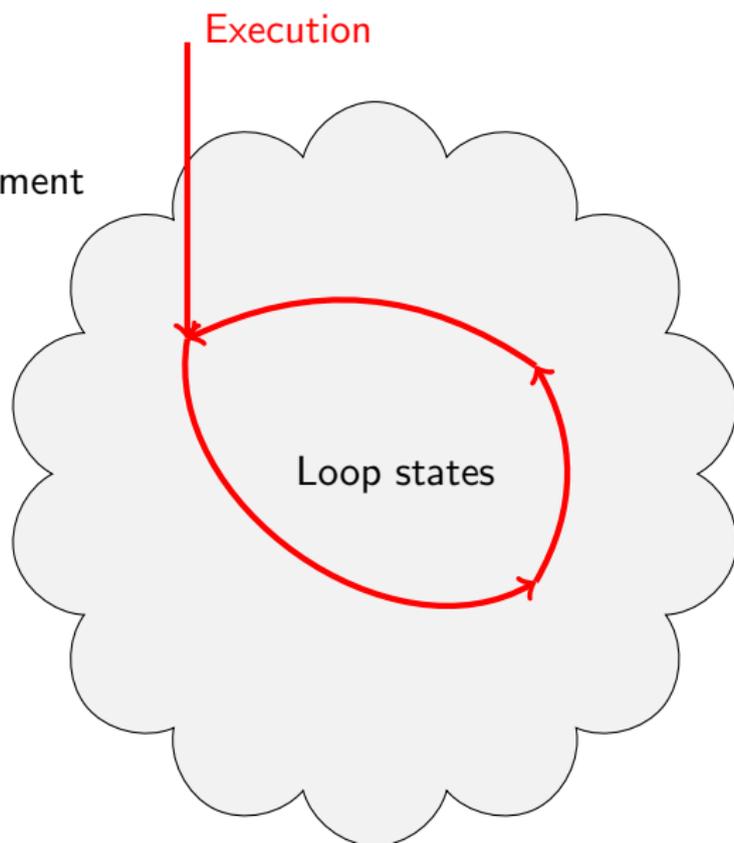
# Termination by iterative strengthening

Find counterexample  
then strengthen argument



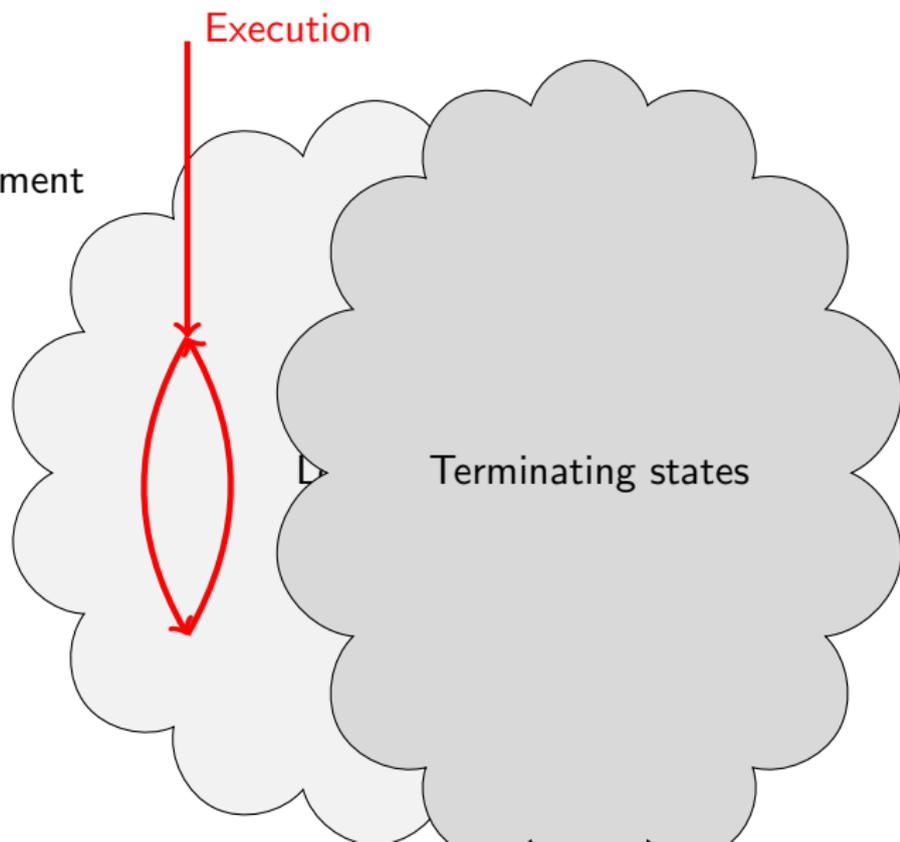
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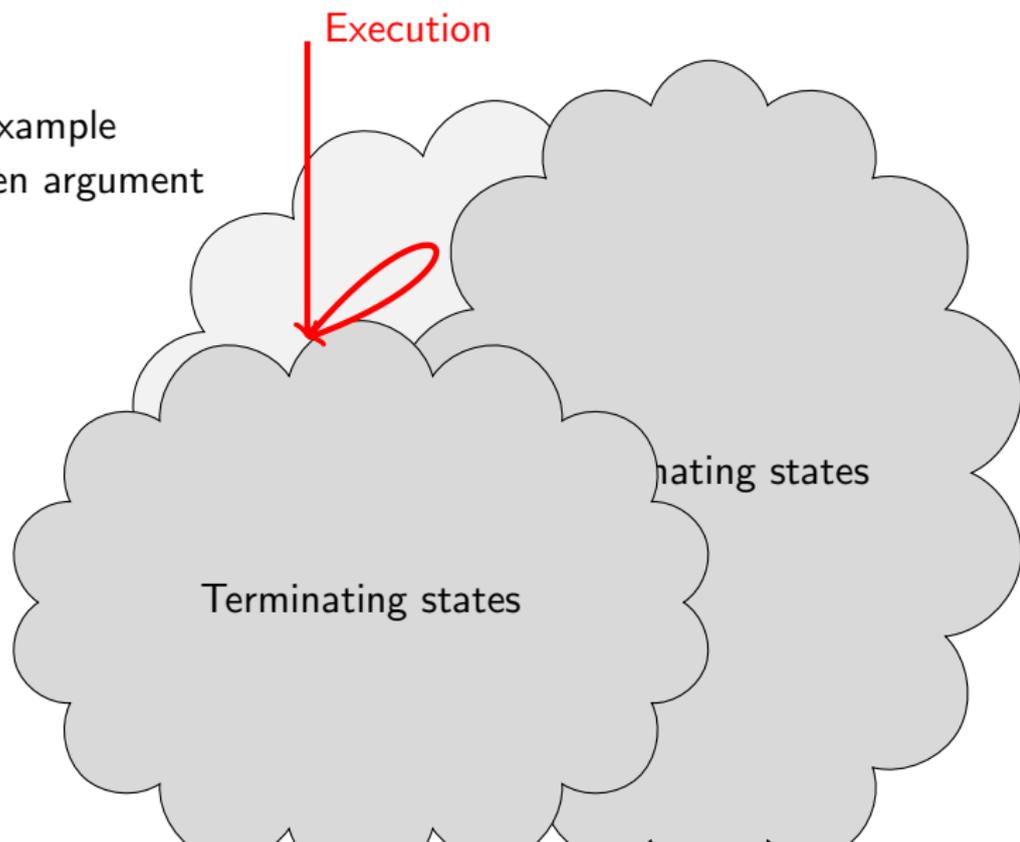
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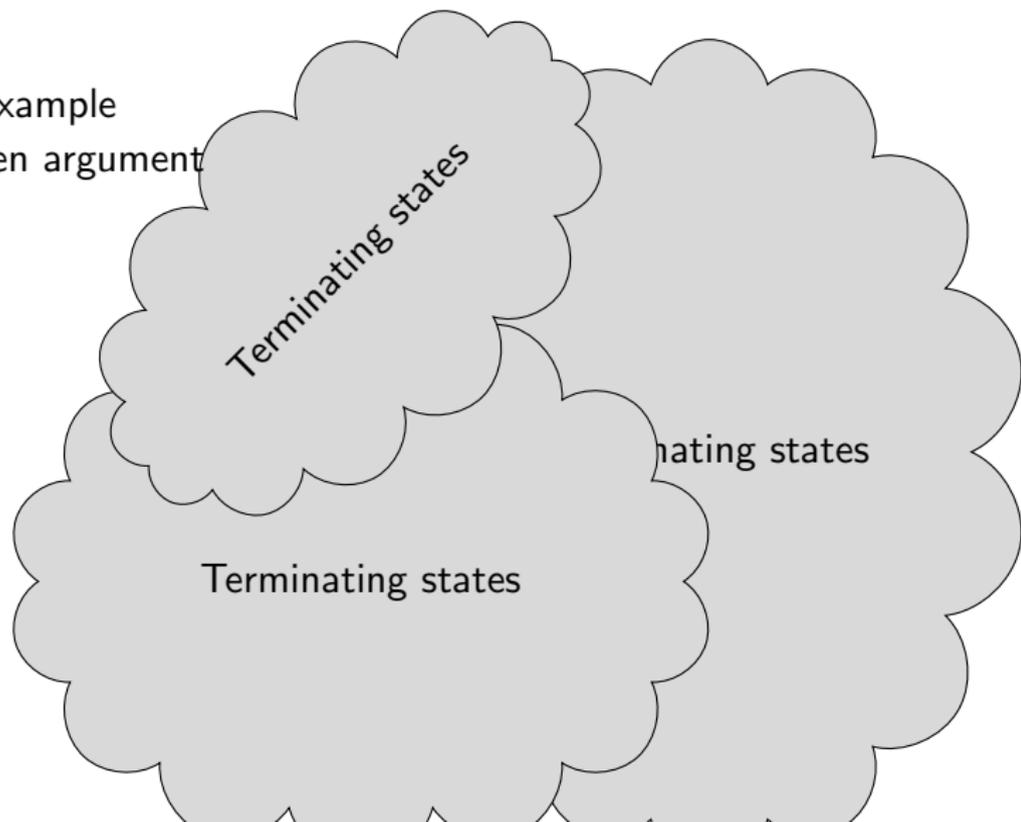
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## Termination by iterative strengthening: Worst case

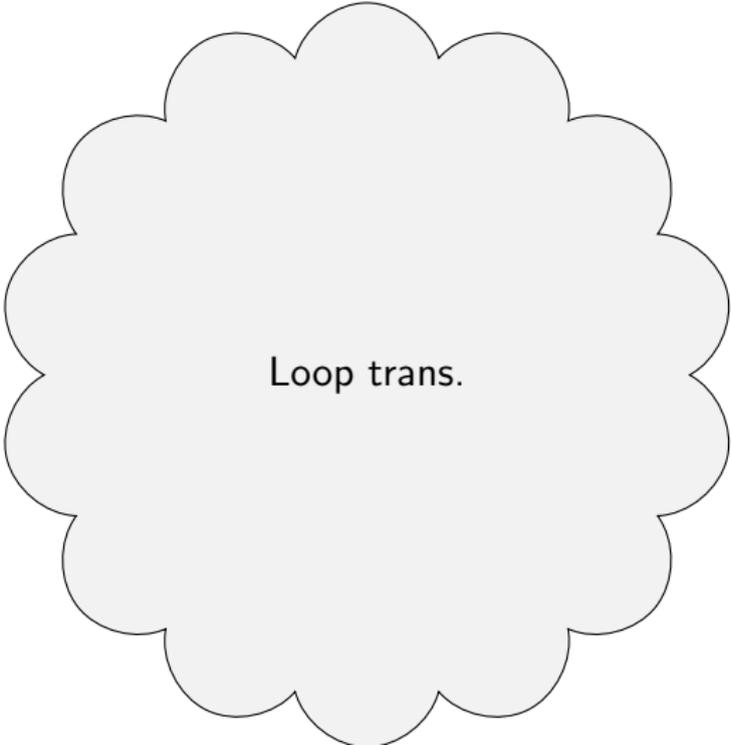
- 1 Safety: Look at everything, then return old sample
- 2 Rank tool: Find **too** specific termination argument
- 3 Safety: Can't prove generality, repeat 1

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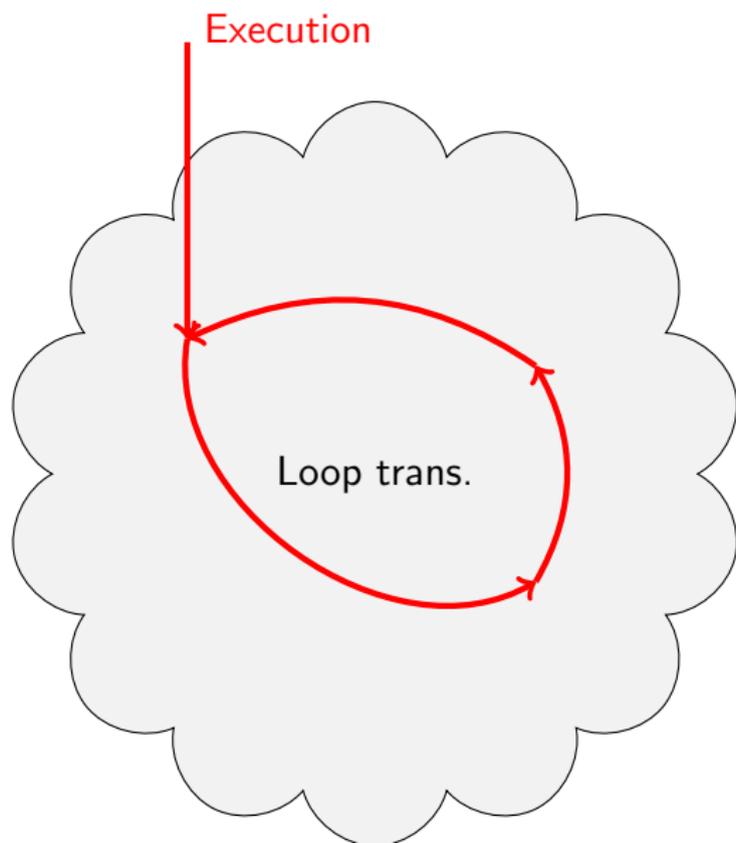


## Termination by iterative simplification



Loop trans.

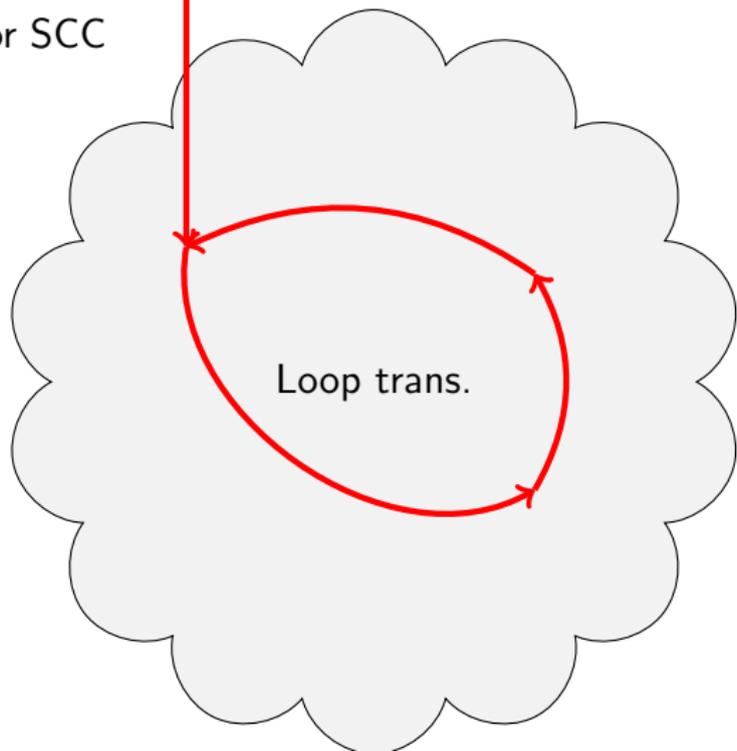
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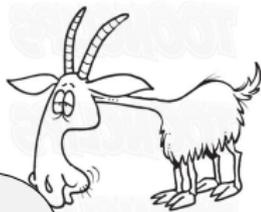
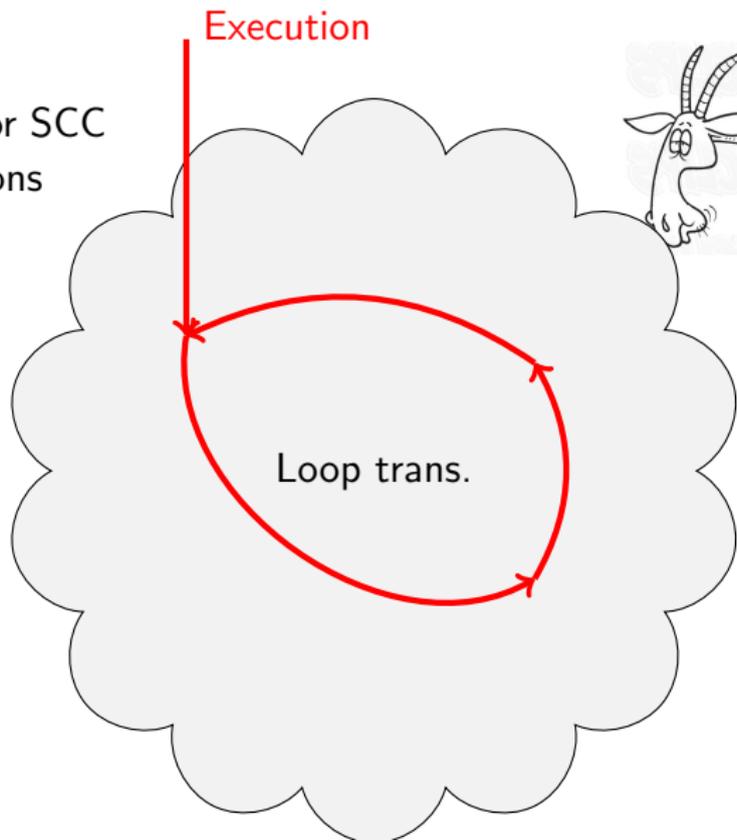
Execution

Find rank function for SCC



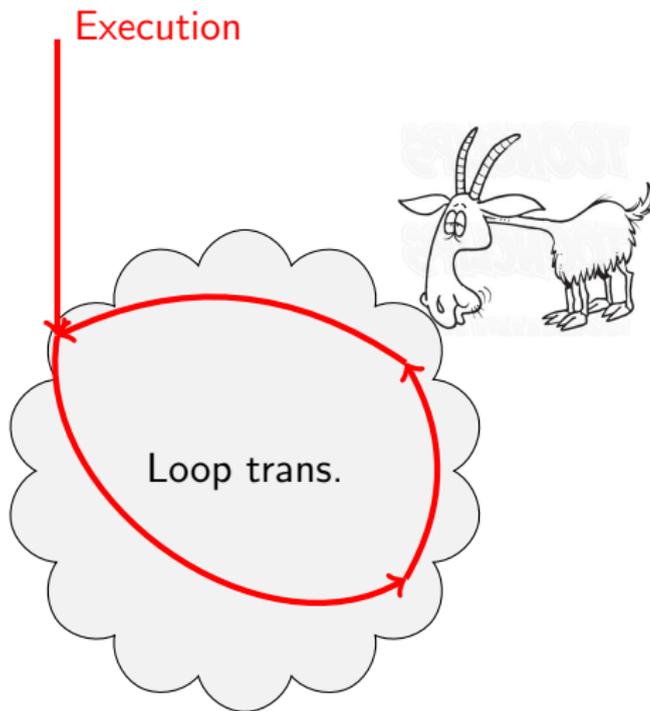
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Find rank function for SCC  
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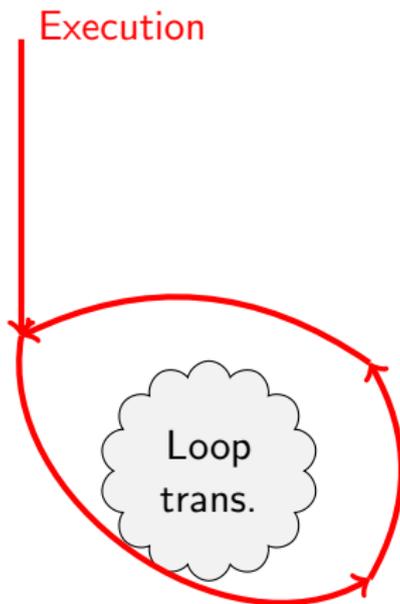
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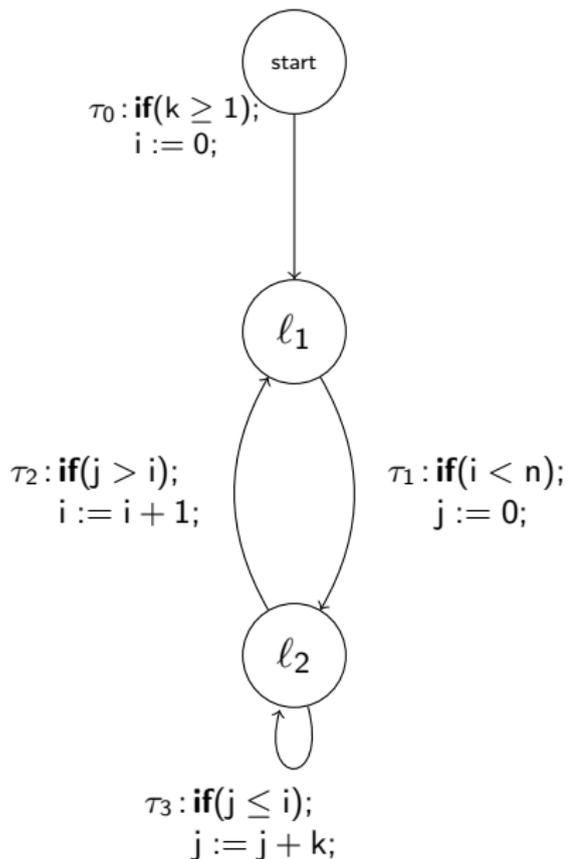
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# Termination by cooperation

- 1 Safety: Provide samples (Counterexamples)
- 2 Rank tool: Find termination argument **in context**
- 3 Rank tool: Mark definitely terminating parts
- 4 Safety: Prove generality for rest, or 1

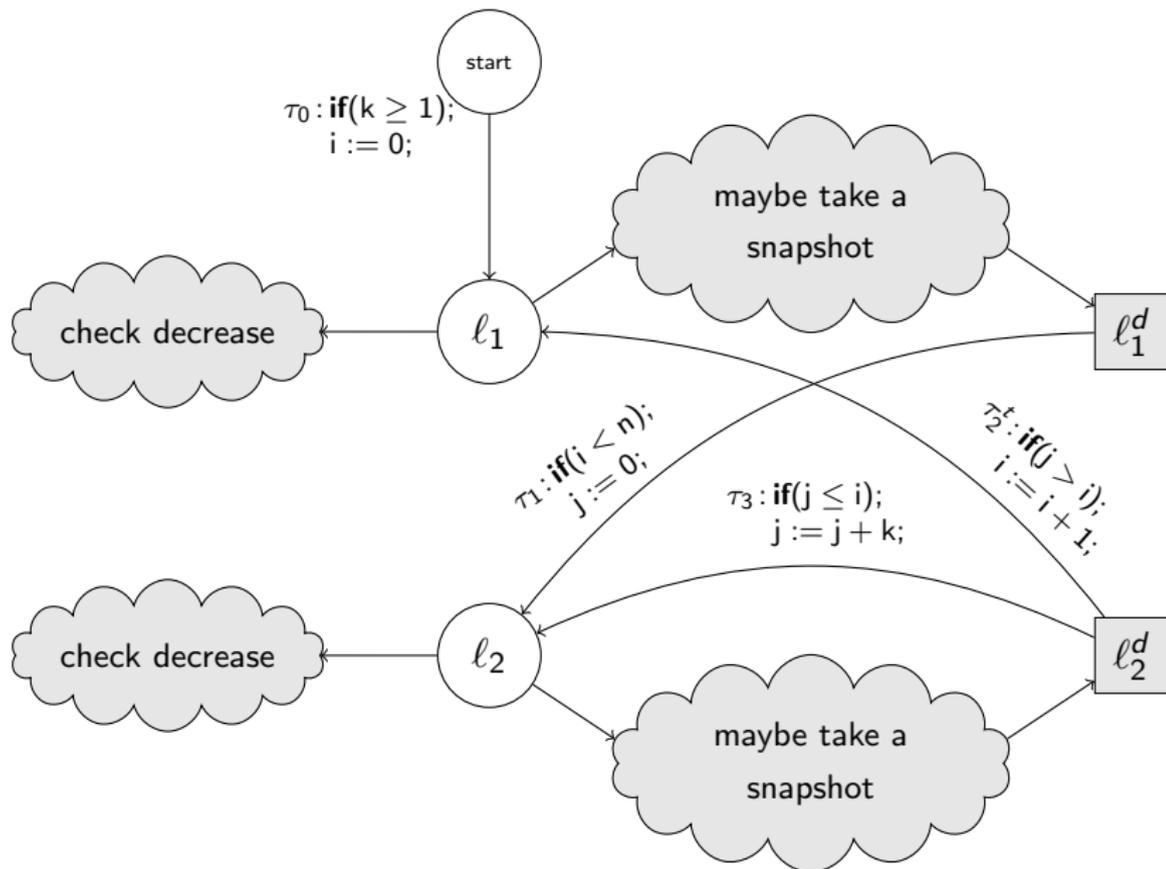
# A nested example



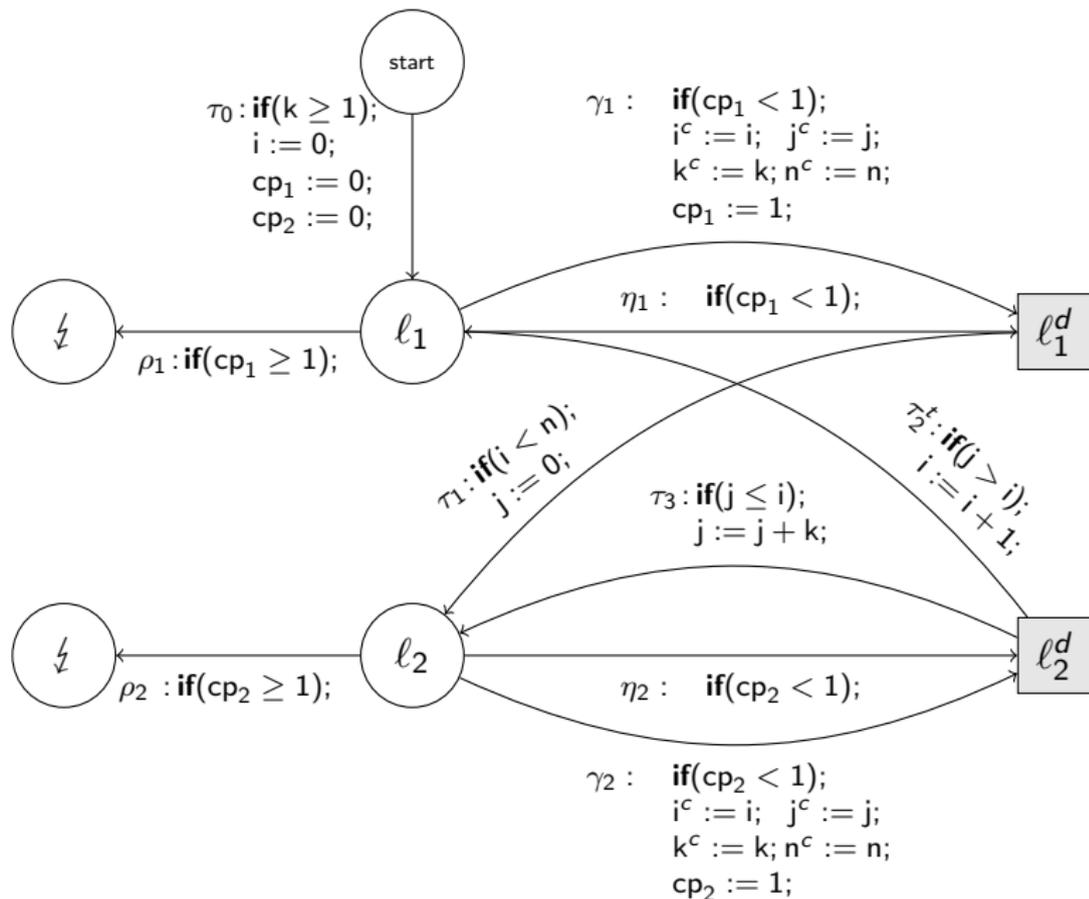
## Example (Source)

```
if  $k \geq 1$  then  
   $i := 0$ ;  
   $l_1$  while  $i < n$  do  
     $j := 0$ ;  
     $l_2$  while  $j \leq i$  do  
       $j := j + k$ ;  
    done  
     $i := i + 1$ ;  
  done  
fi
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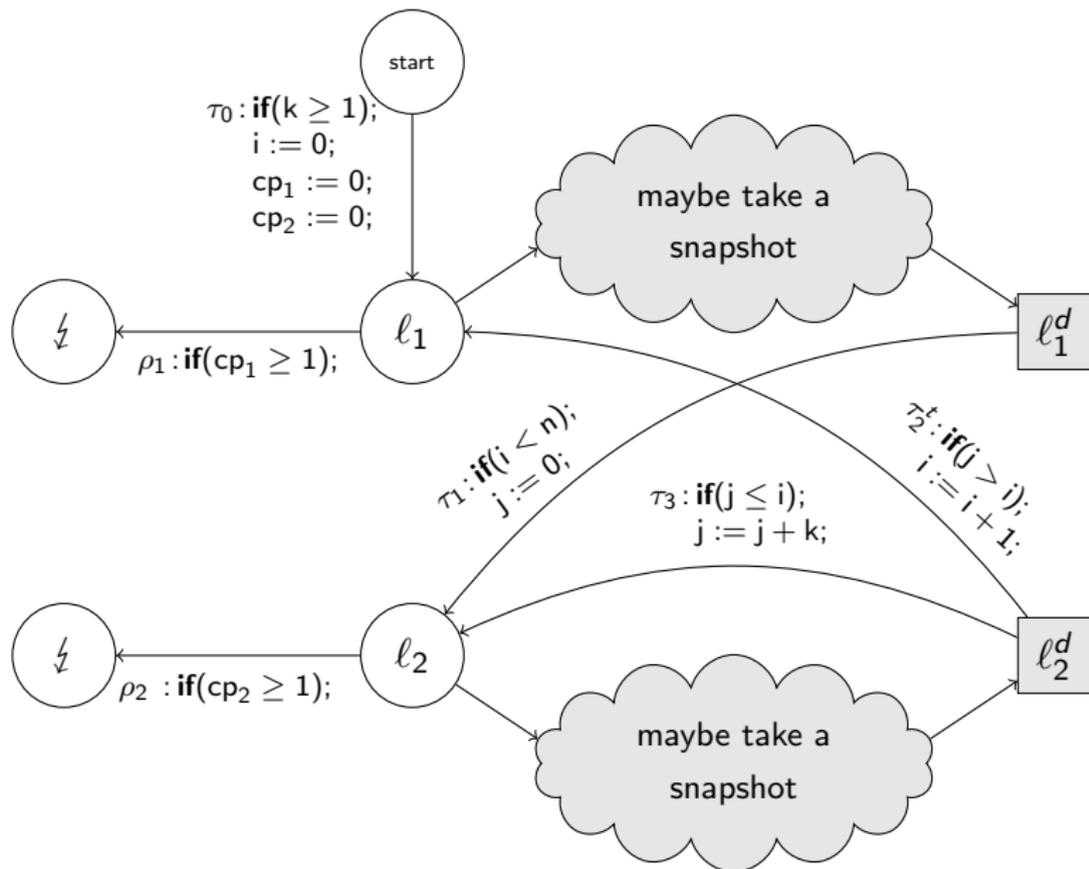
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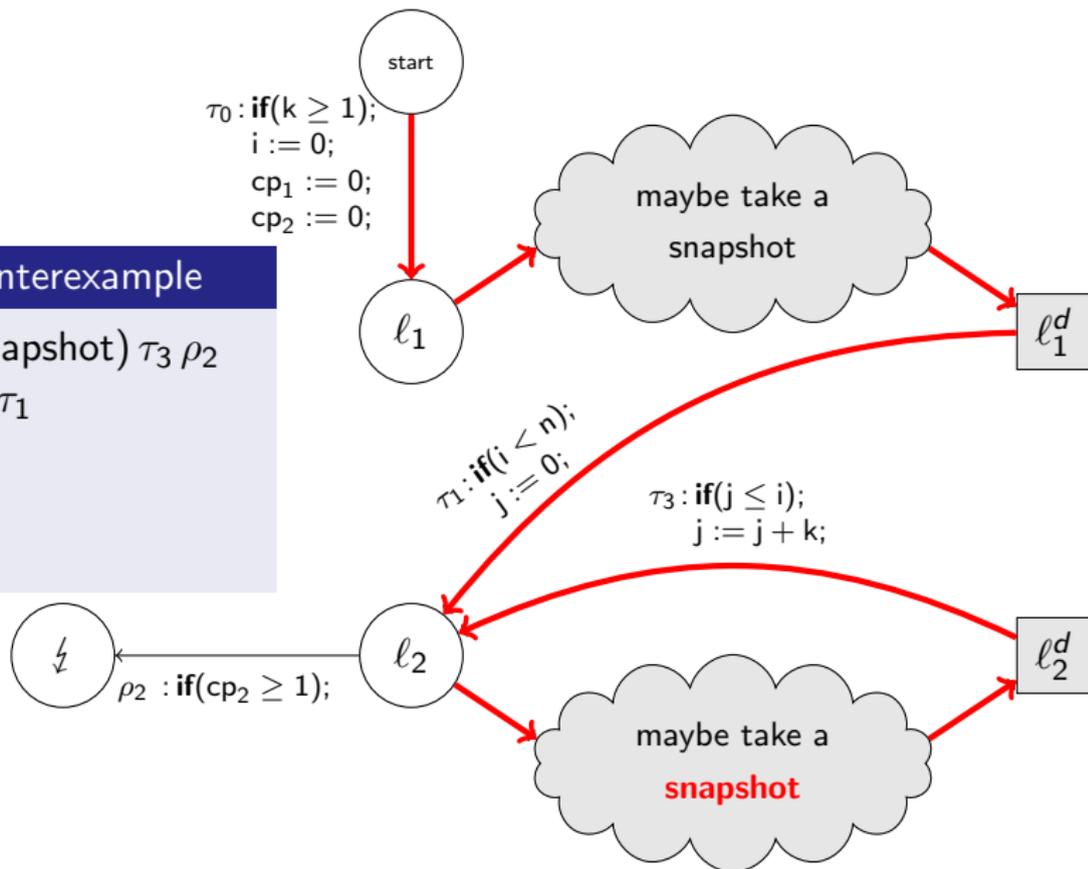
# A nested example: Strengthening

## First Counterexample

$\tau_0 \tau_1$  (snapshot)  $\tau_3 \rho_2$

Stem:  $\tau_0 \tau_1$

Loop:  $\tau_3$



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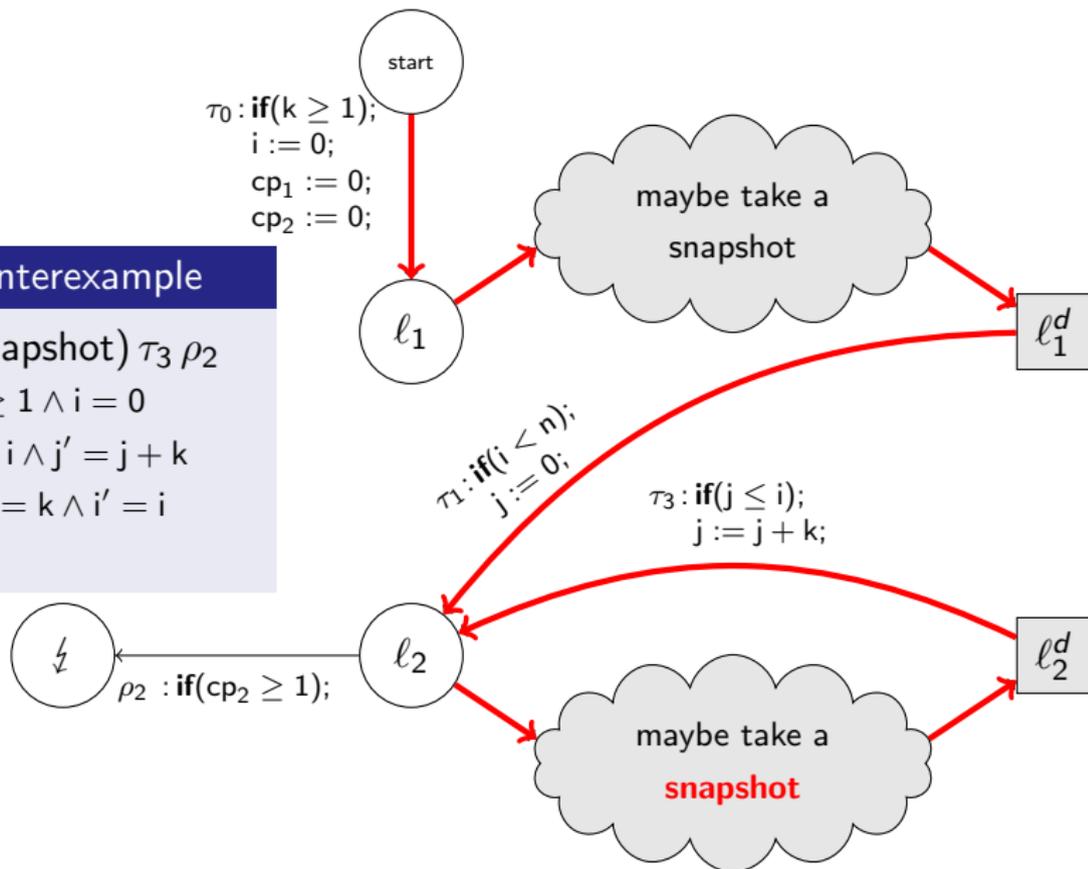
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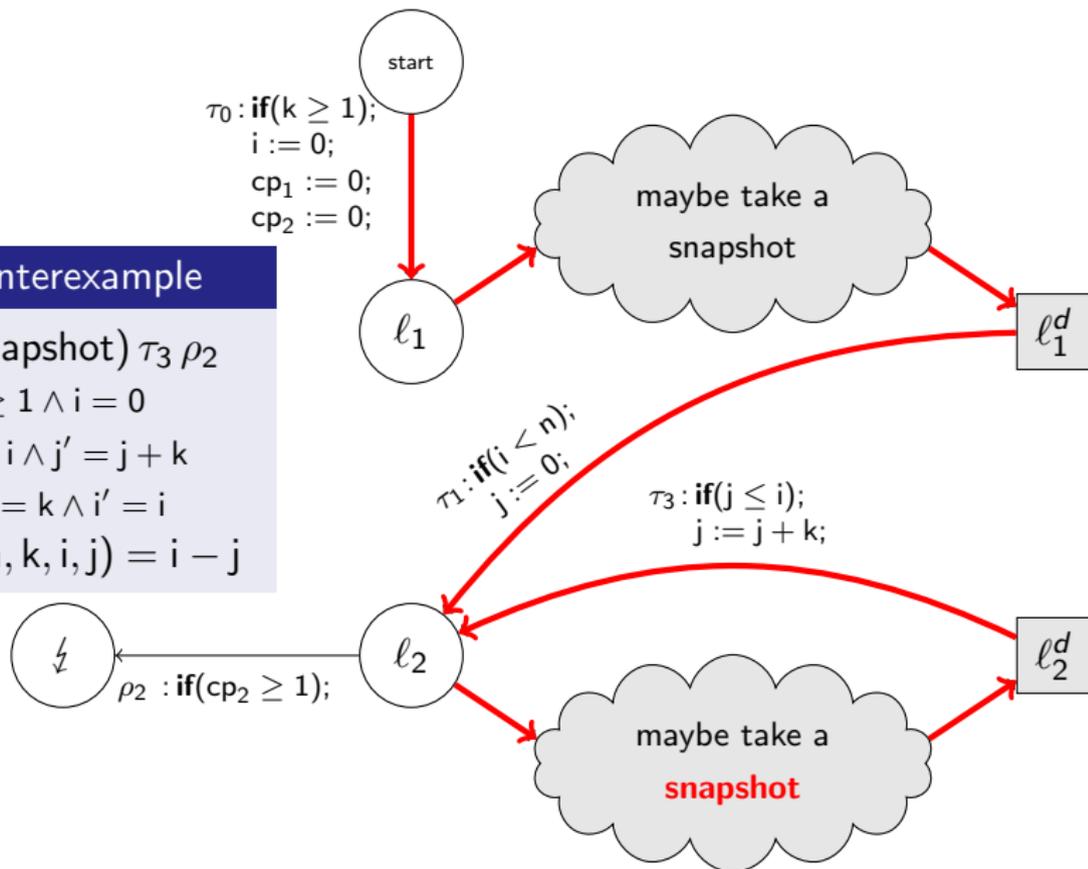
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RF:  $f_{l_2}(n, k, i, j) = i - j$



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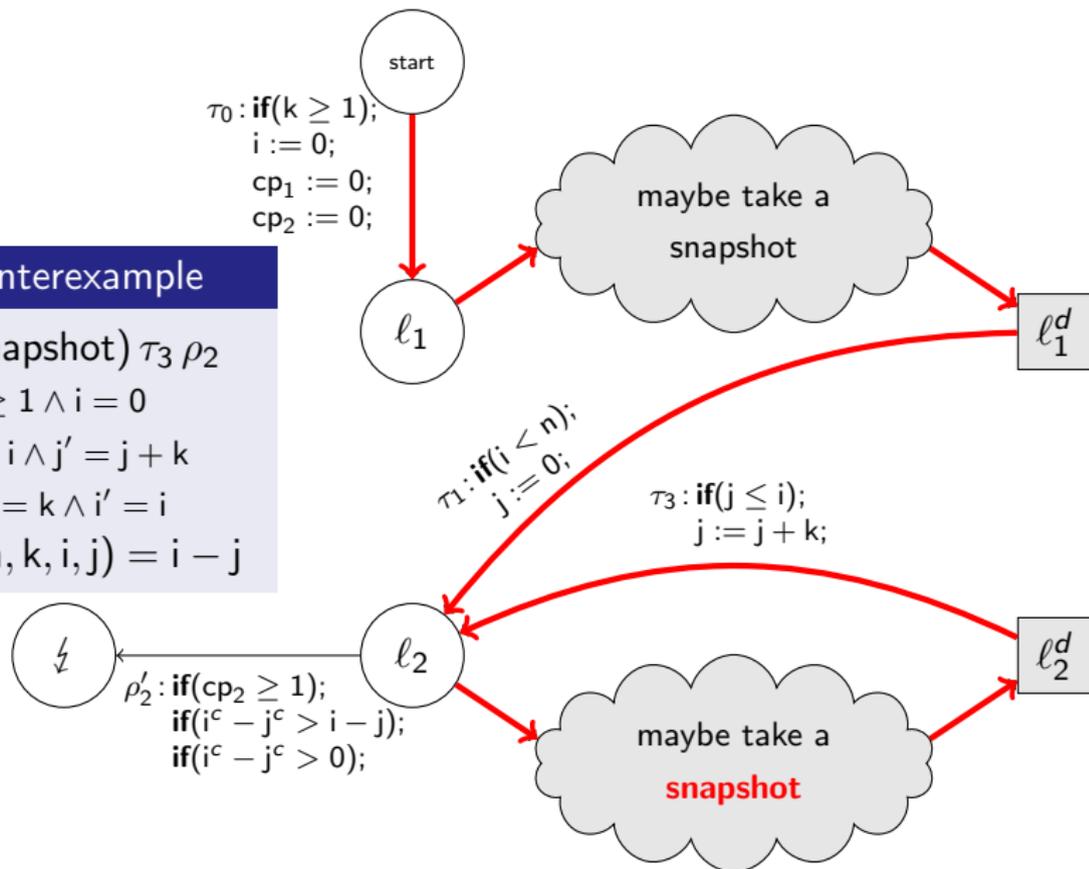
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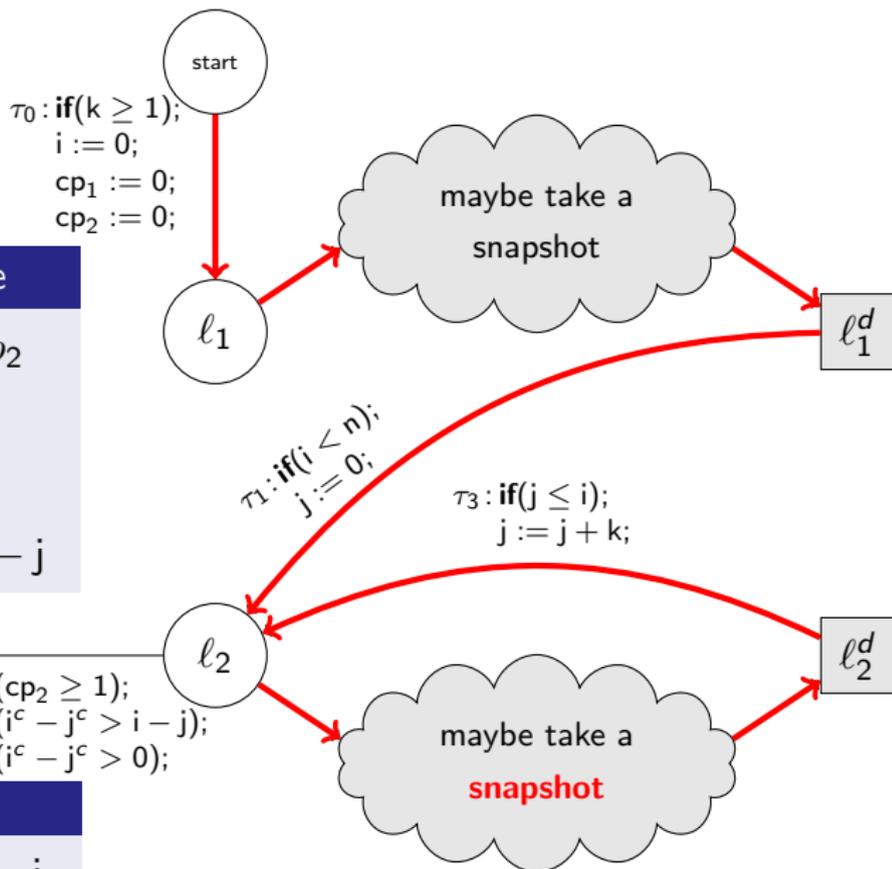
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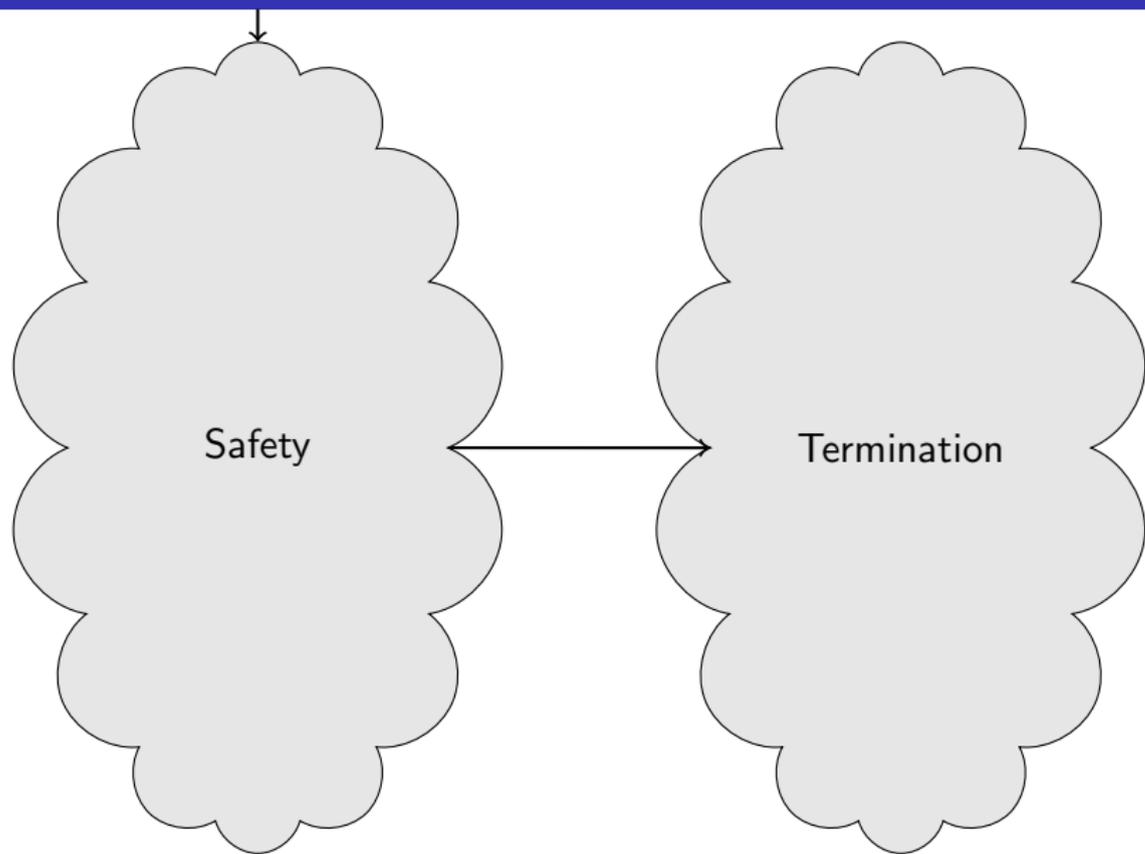
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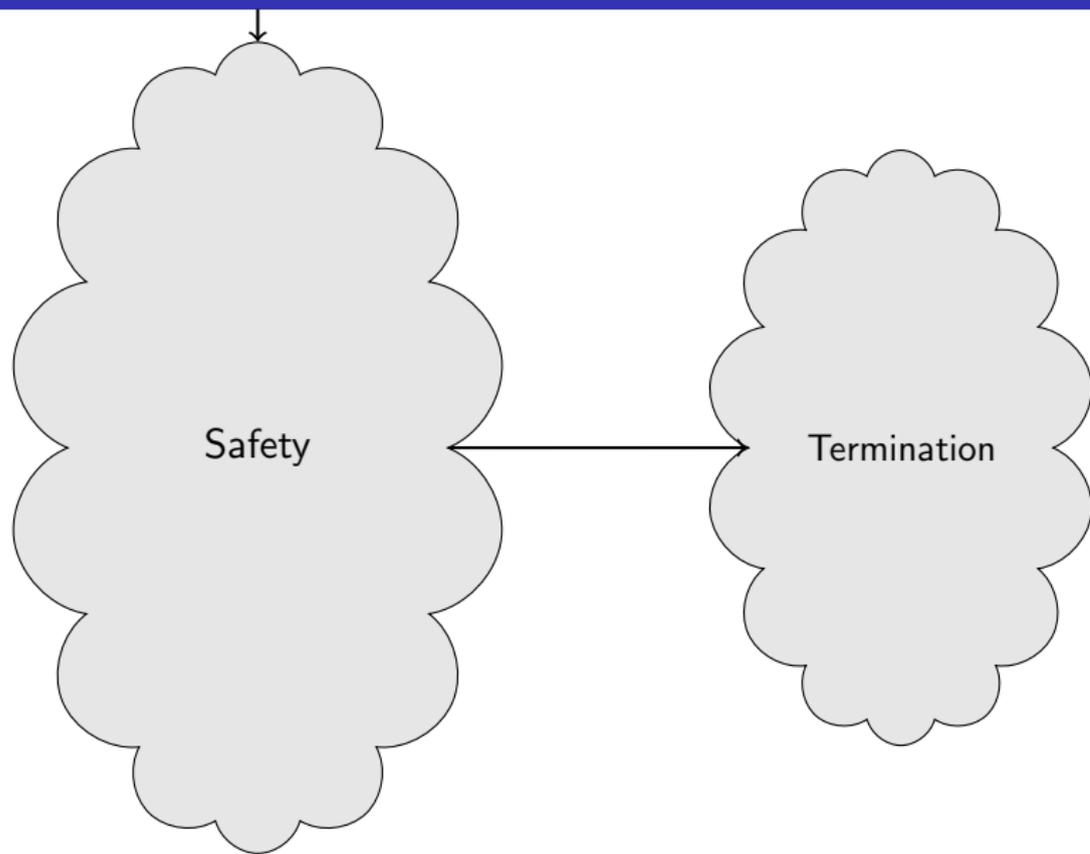
## Alternative RF

RF:  $f'_{l_2}(n, k, i, j) = 1 - j$

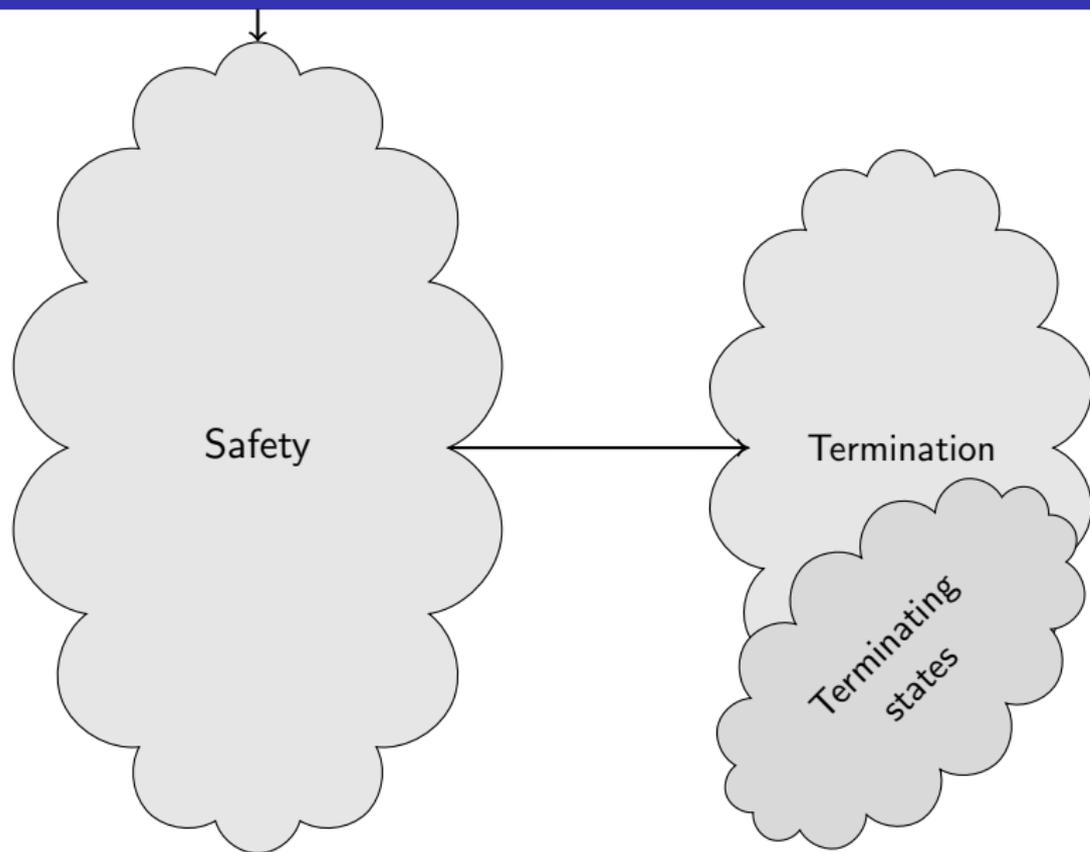
## Cooperation: High-level view



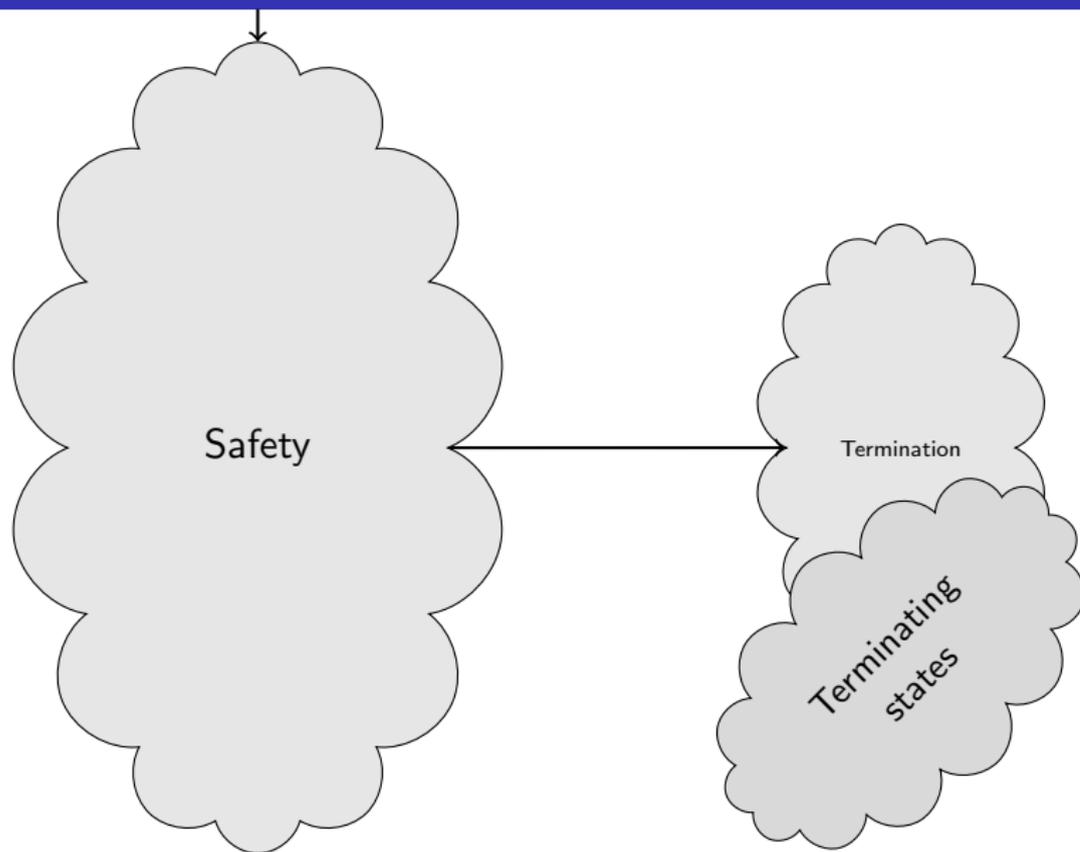
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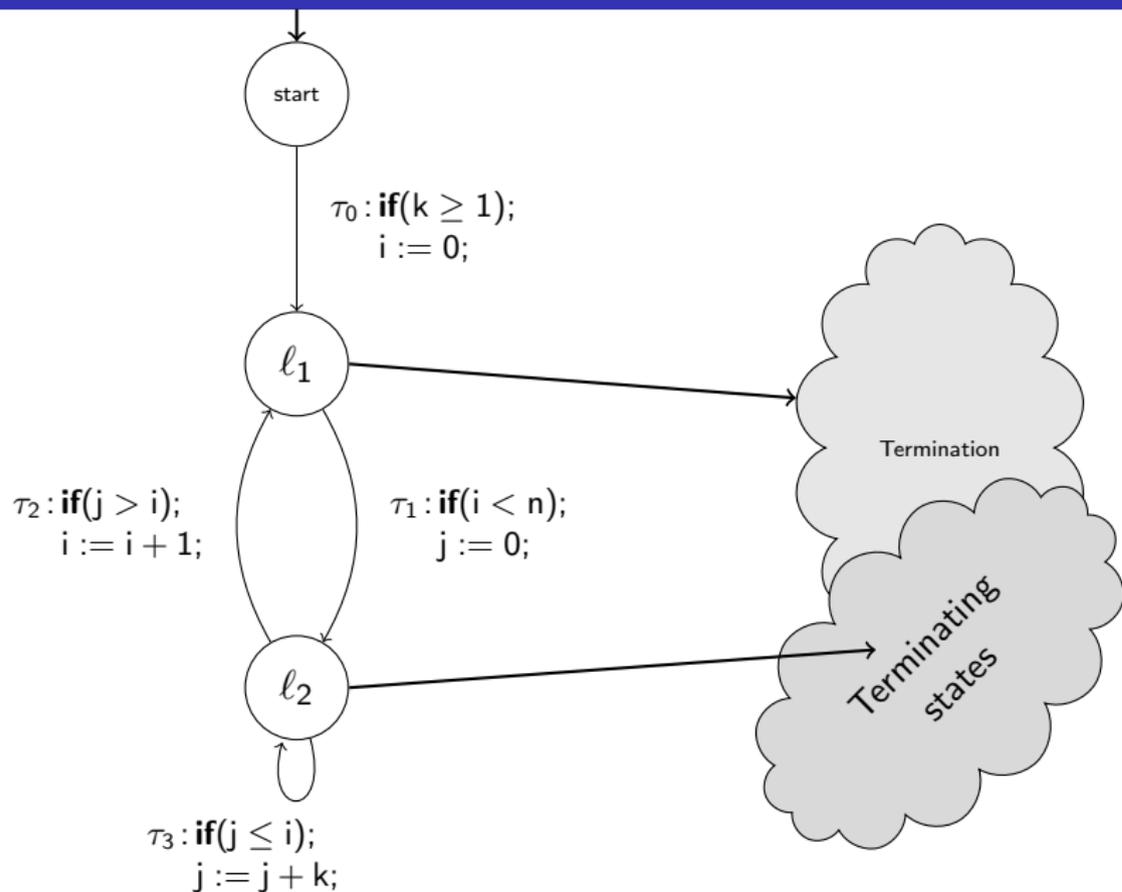
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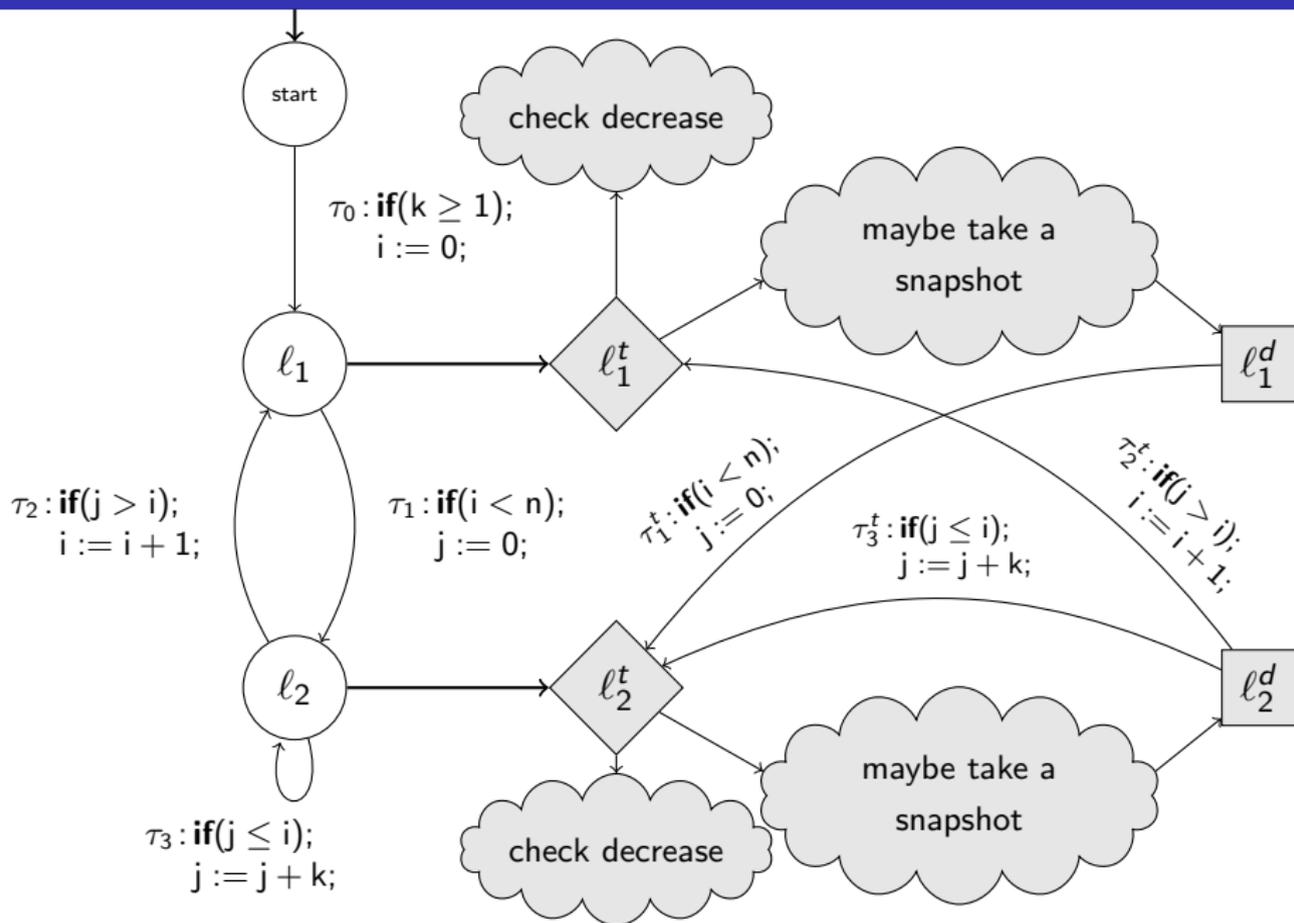
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Approach:

- Analyze whole SCC, not counterexample slice

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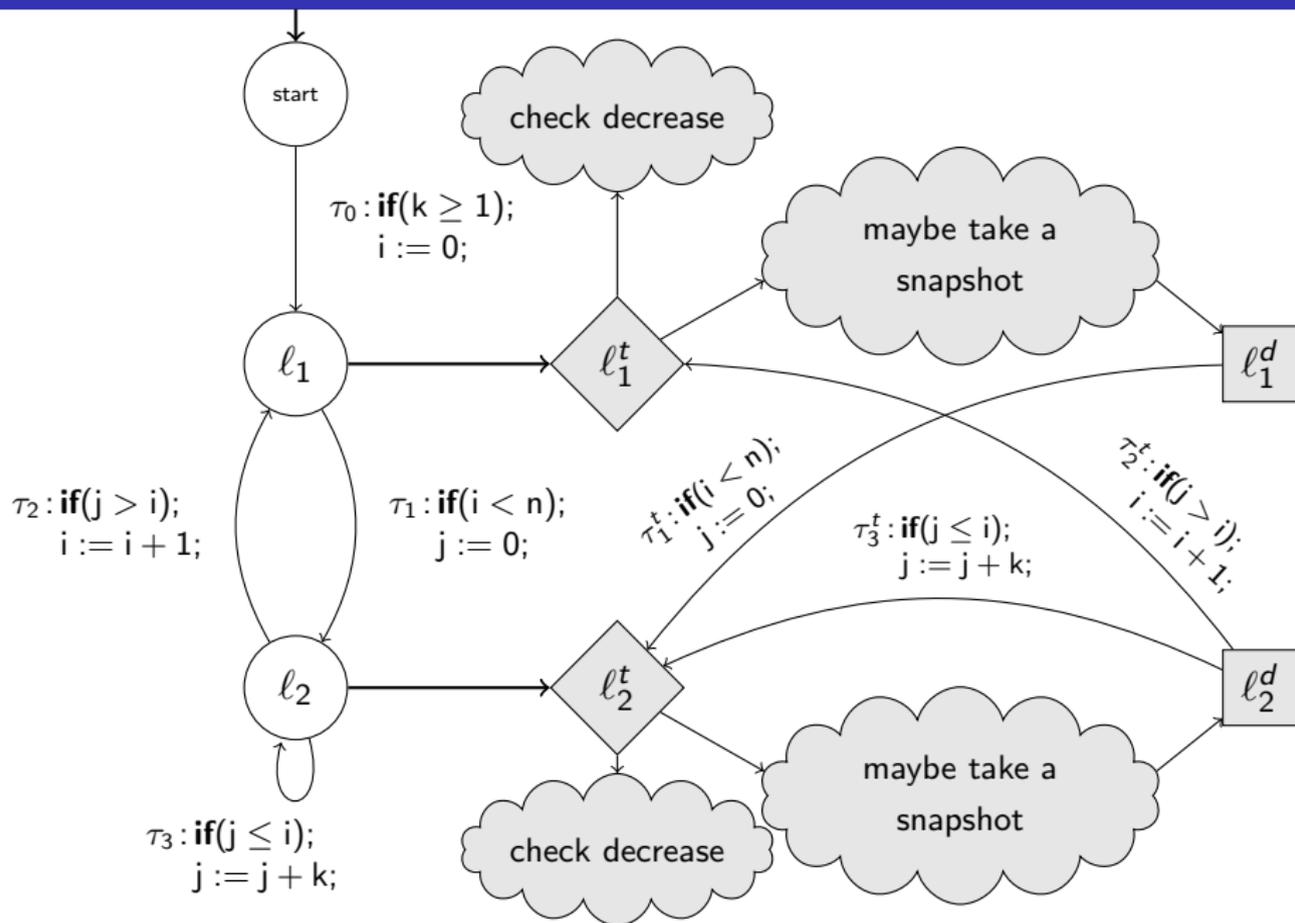
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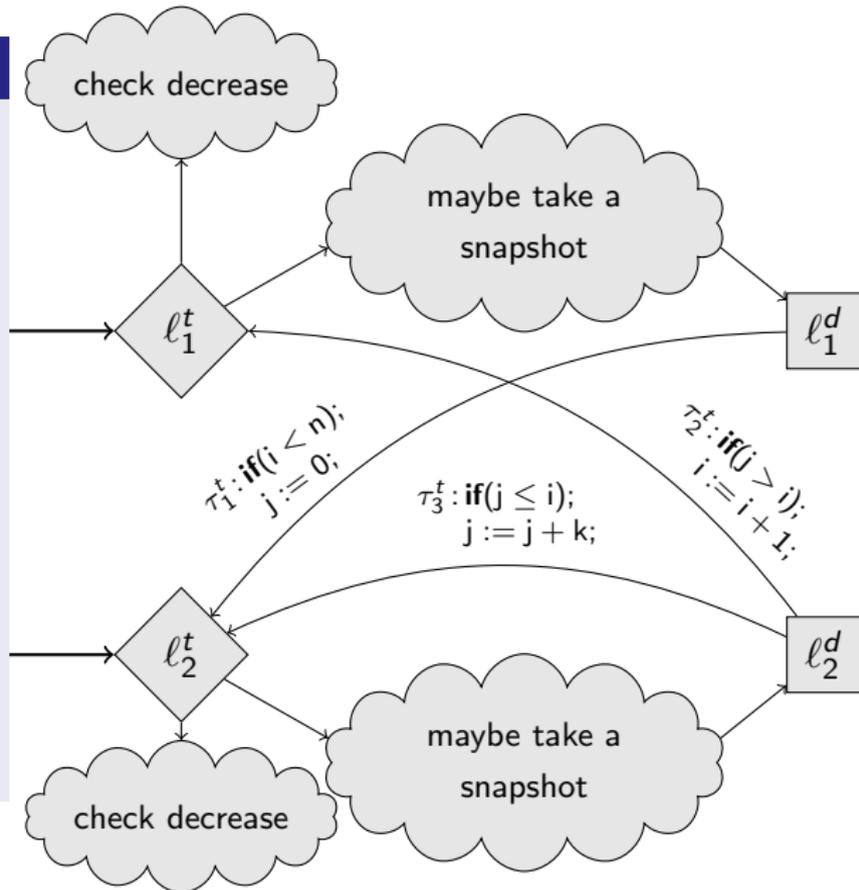
- Analyze whole SCC, not counterexample slice
- Remove transitions after proof

# Cooperation: Simplification



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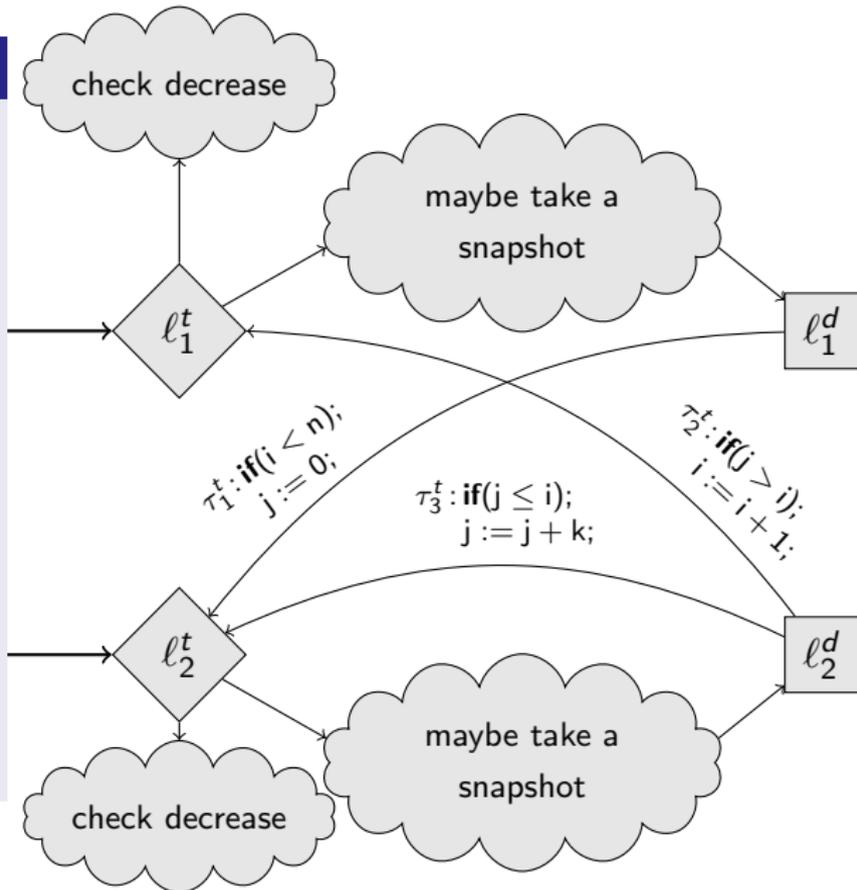


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## Simplification

- 1 Find SCC  $\mathcal{S}$  in termination graph:

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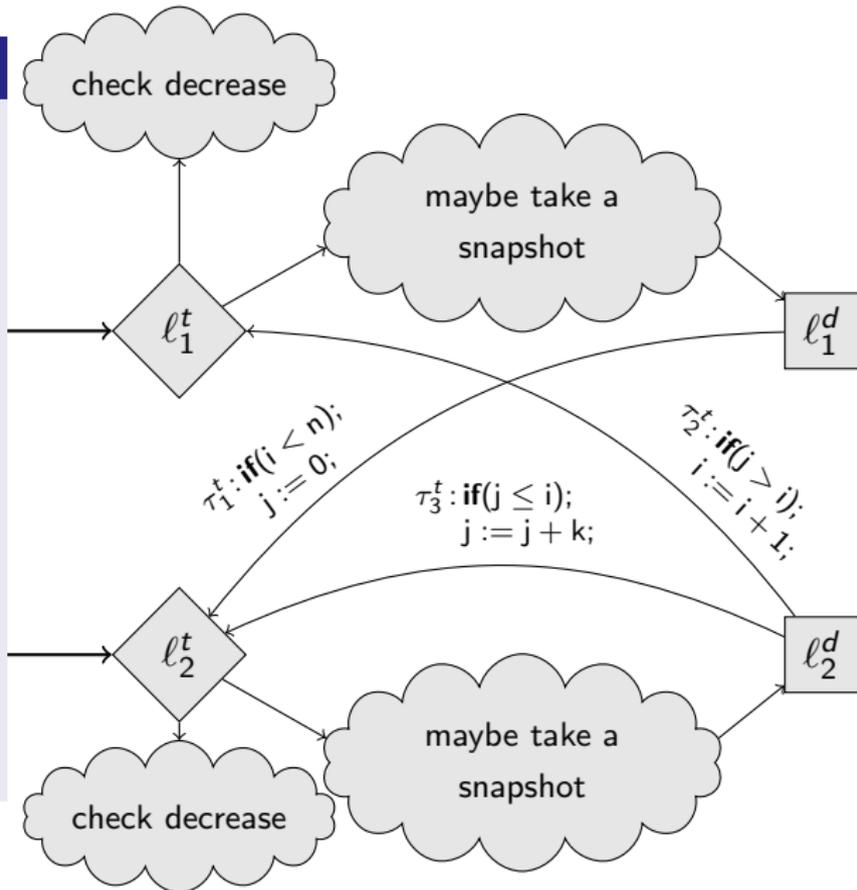
- 2 Find  $\mathcal{S}$ -orienting RF:

$$f_{l_1^t}^1(i, j, k, n) = n - i + 1$$

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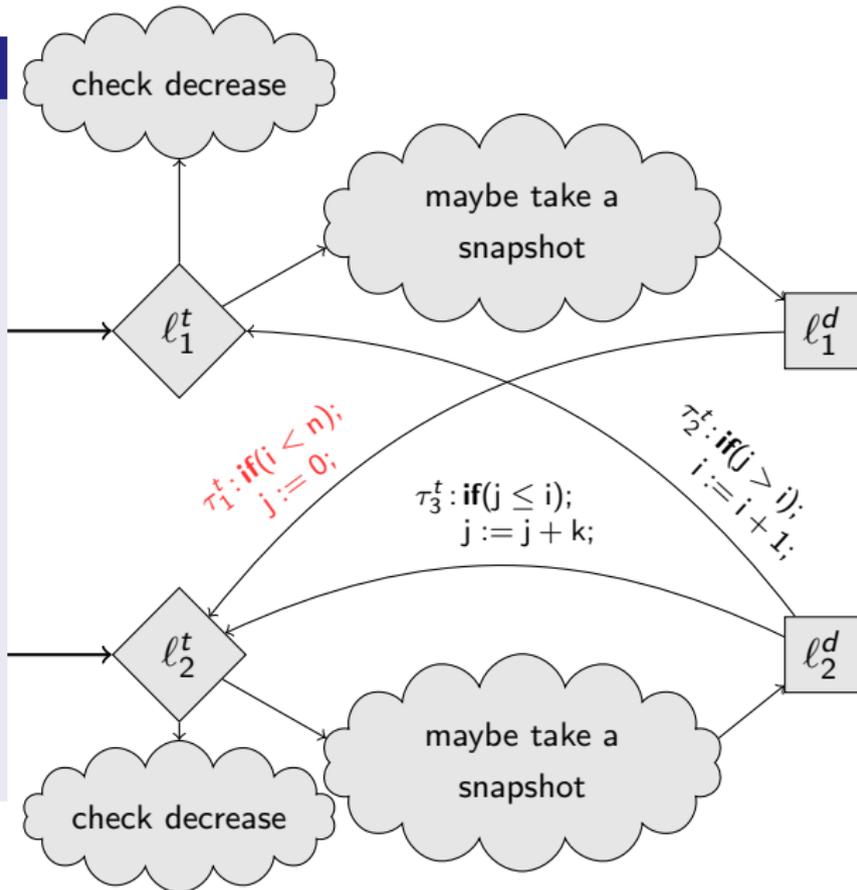
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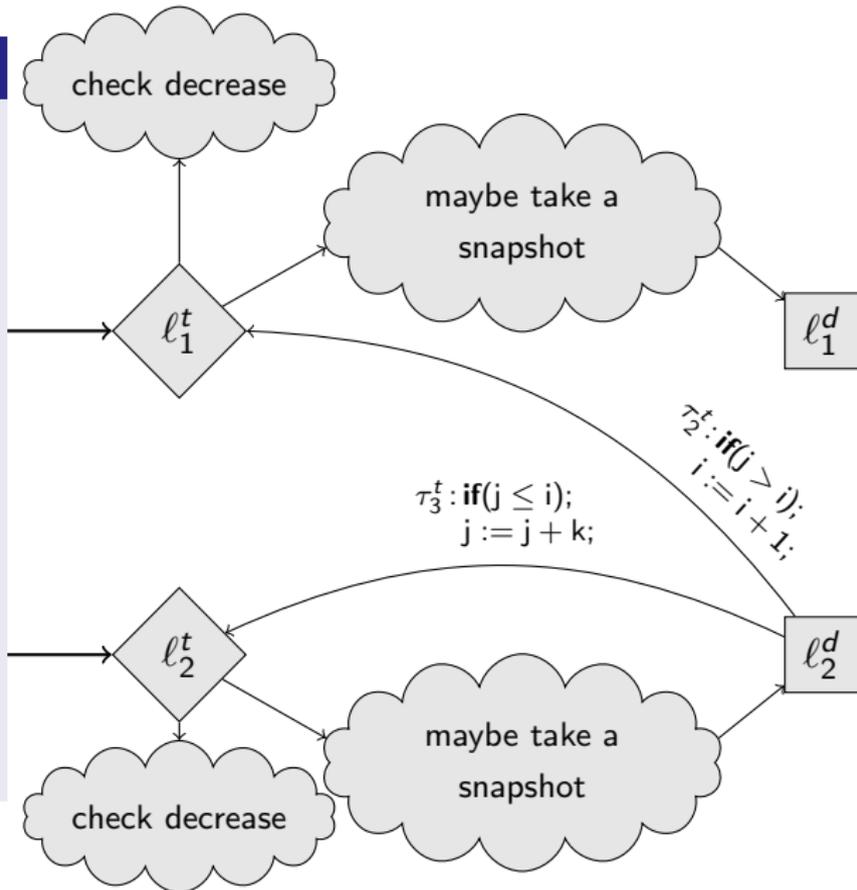
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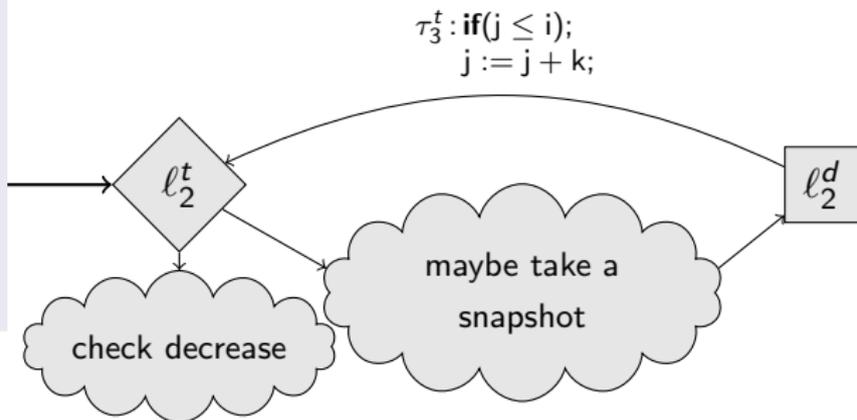
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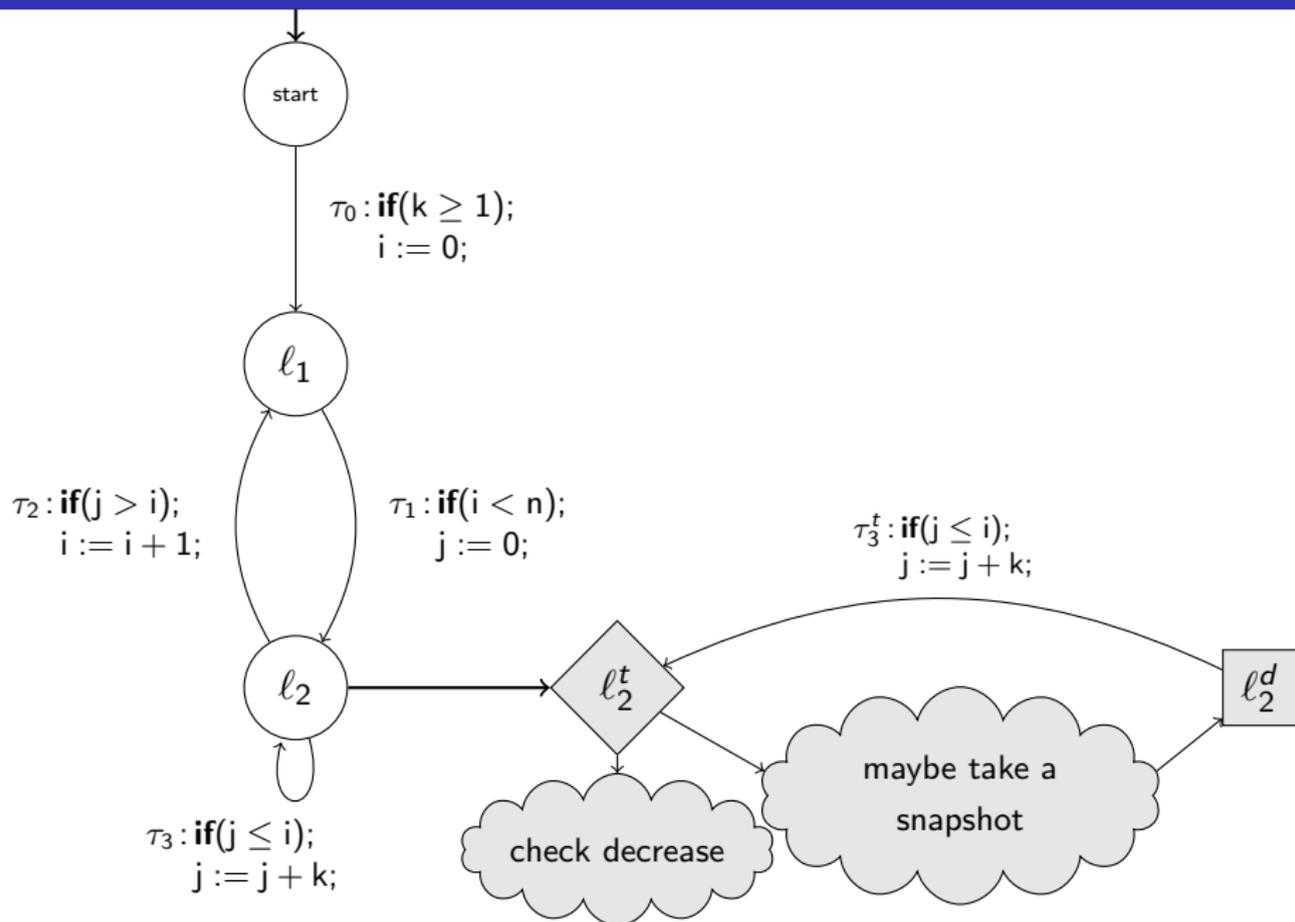
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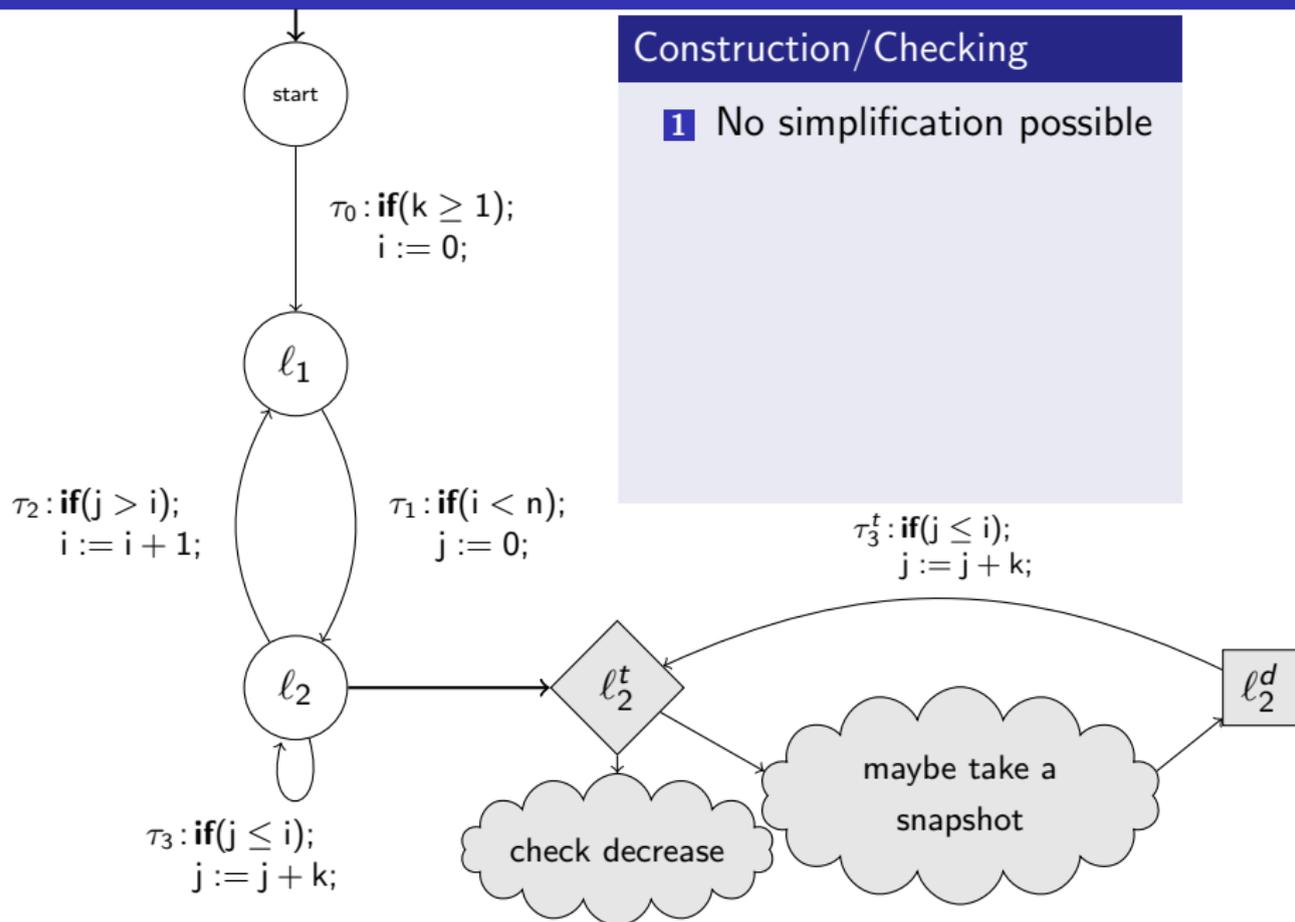
- 3 Delete decr./bounded
- 4 Clean up



# Cooperation



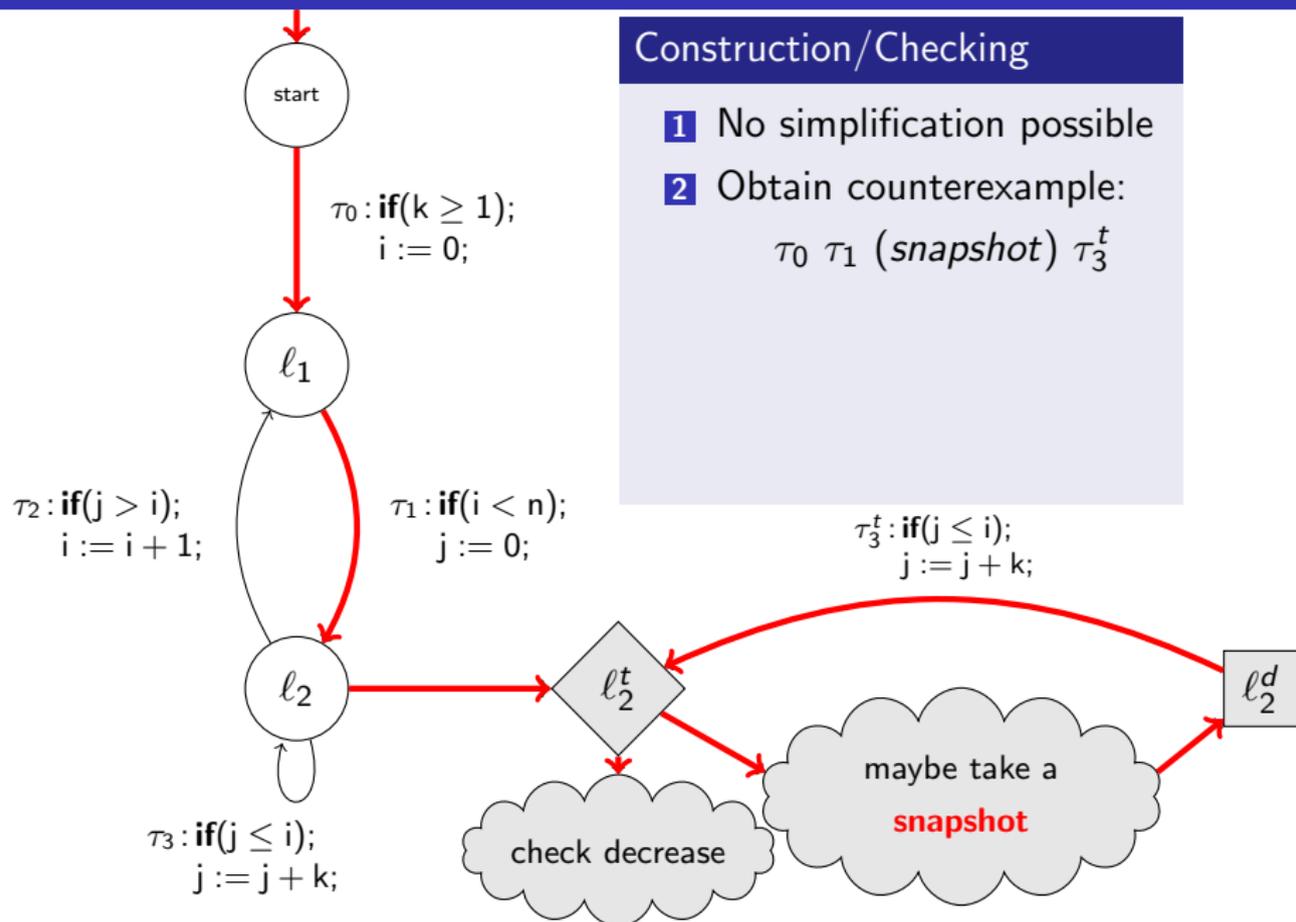
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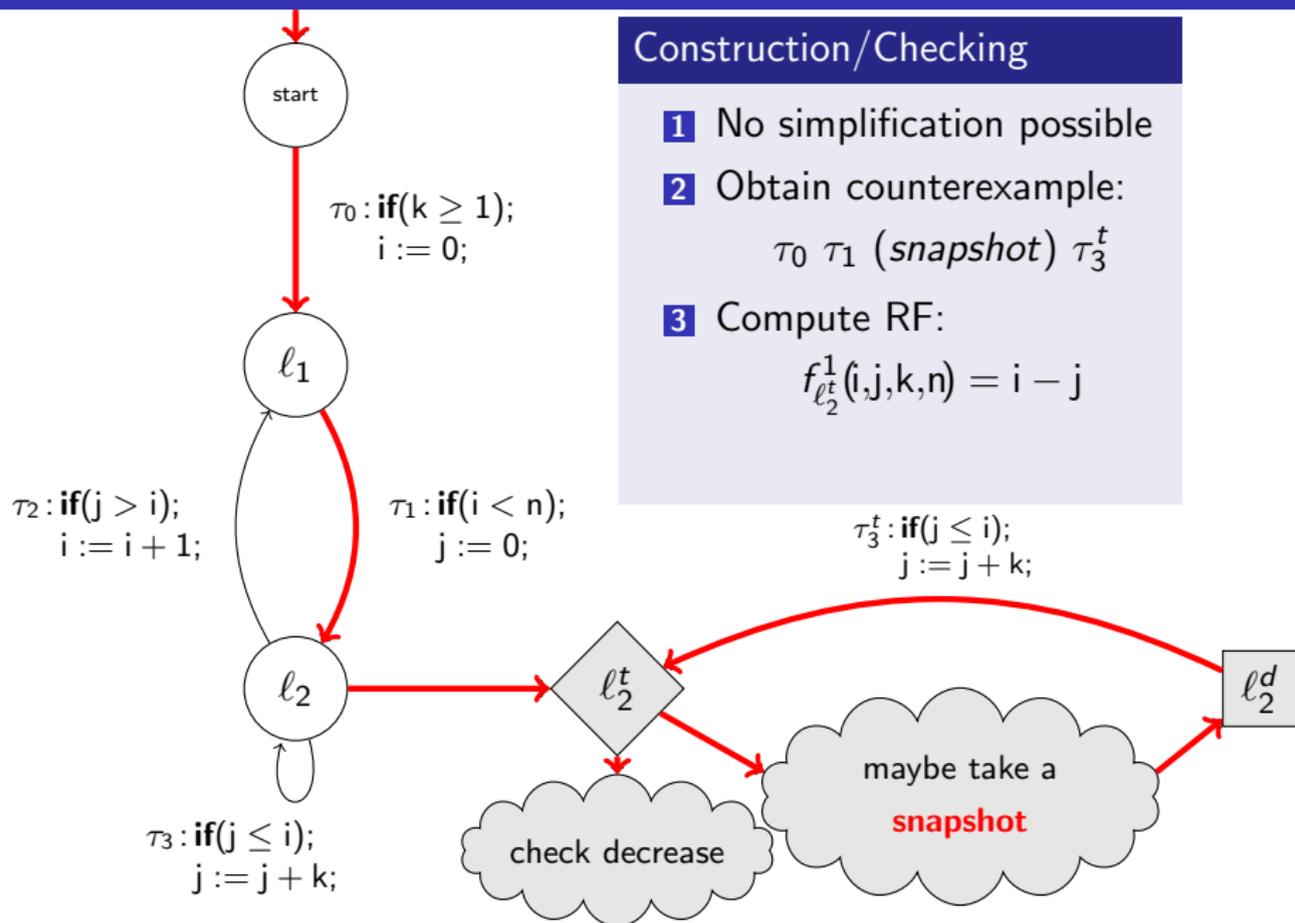
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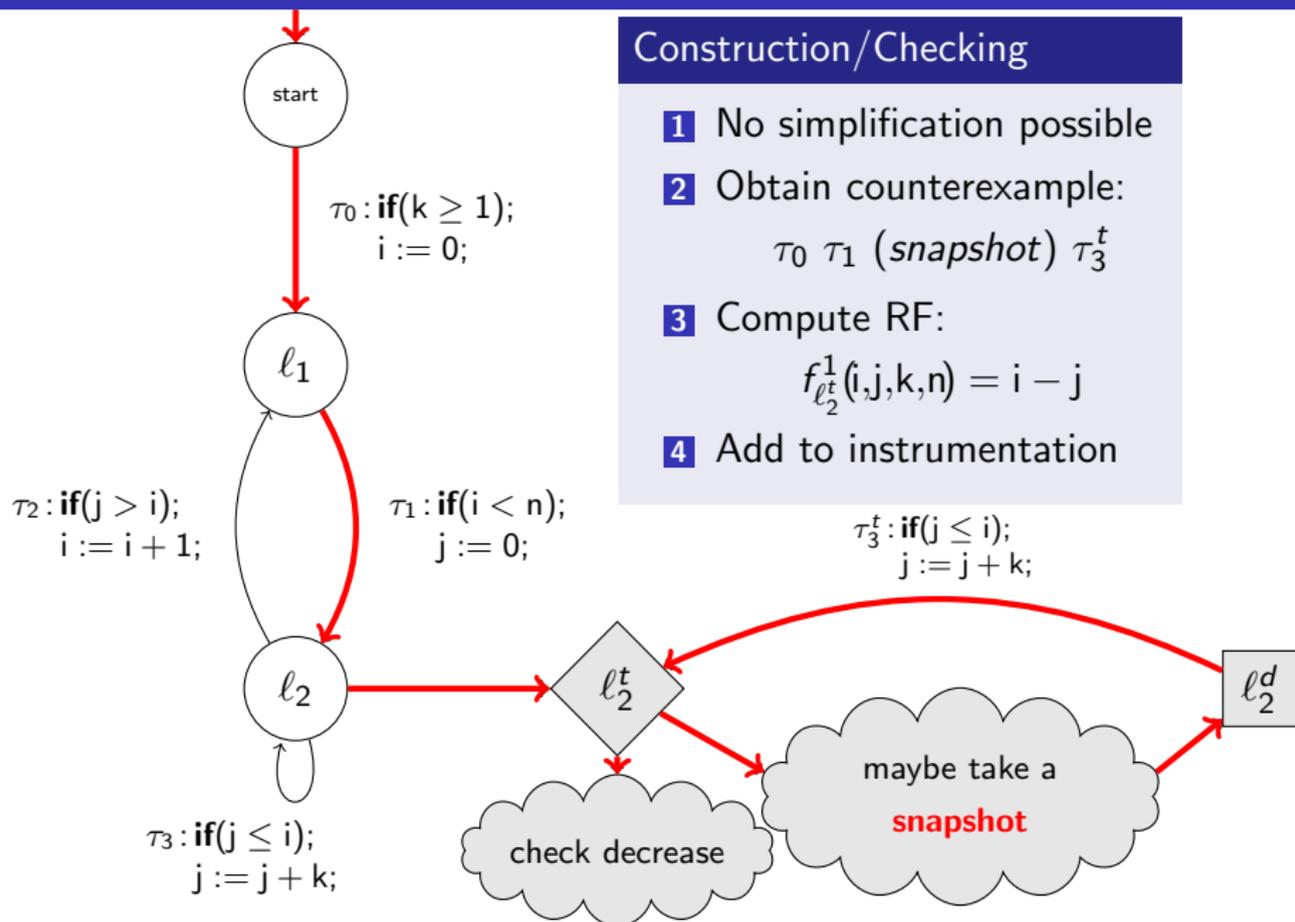
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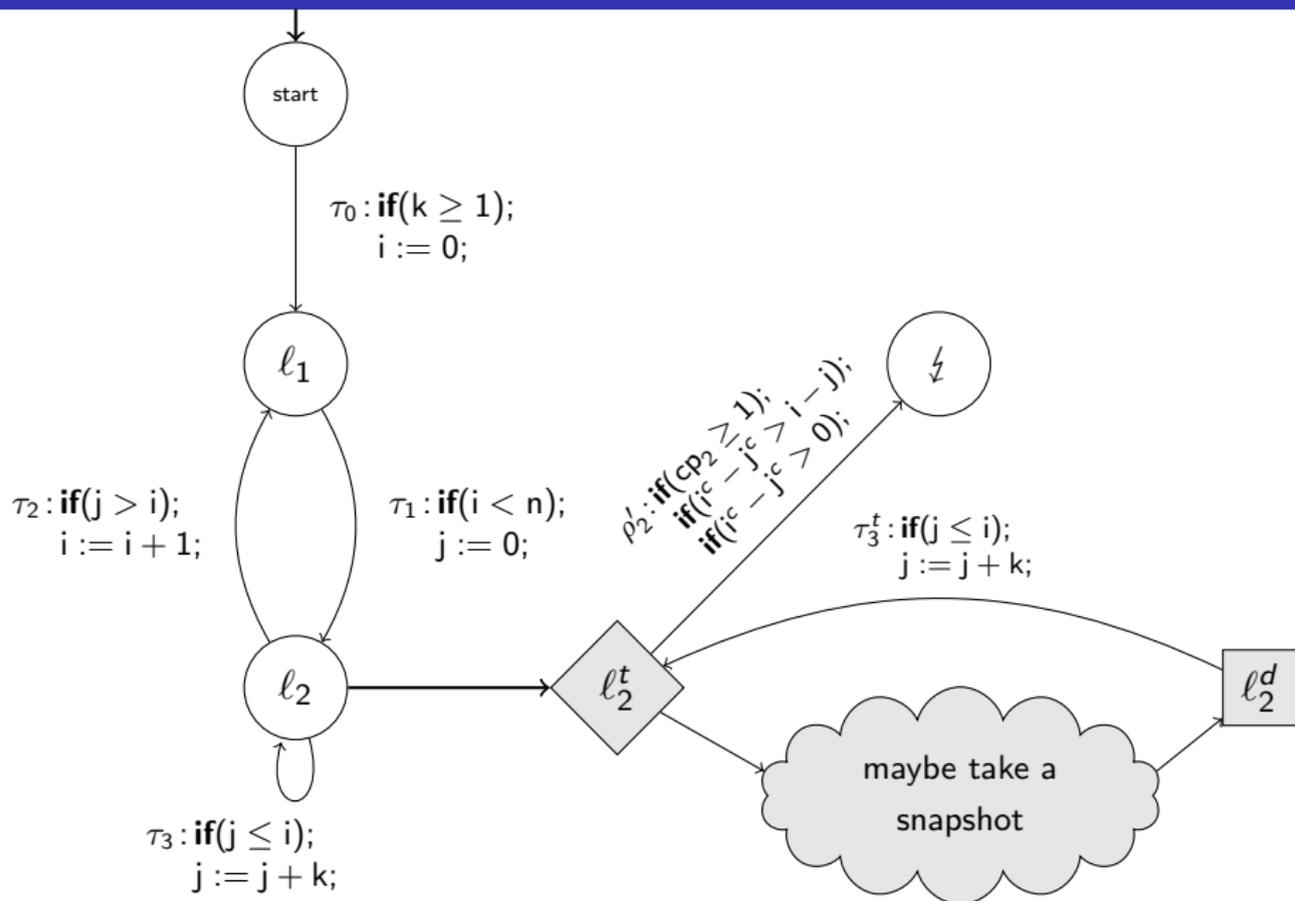


## Construction/Checking

- 1 No simplification possible
- 2 Obtain counterexample:  
 $\tau_0 \tau_1 (\text{snapshot}) \tau_3^t$
- 3 Compute RF:  
 $f_{l_2^t}^1(i, j, k, n) = i - j$
- 4 Add to instrumentation

$\tau_3^t: \mathbf{if}(j \leq i);$   
 $j := j + k;$

# Cooperation: Invariants



# Cooperation: Evaluation

Evaluated on 449 termination proving benchmarks

260 known terminating, 181 known non-terminating, 8 unknown

Sources: Windows drivers, APACHE, POSTGRESQL, ...

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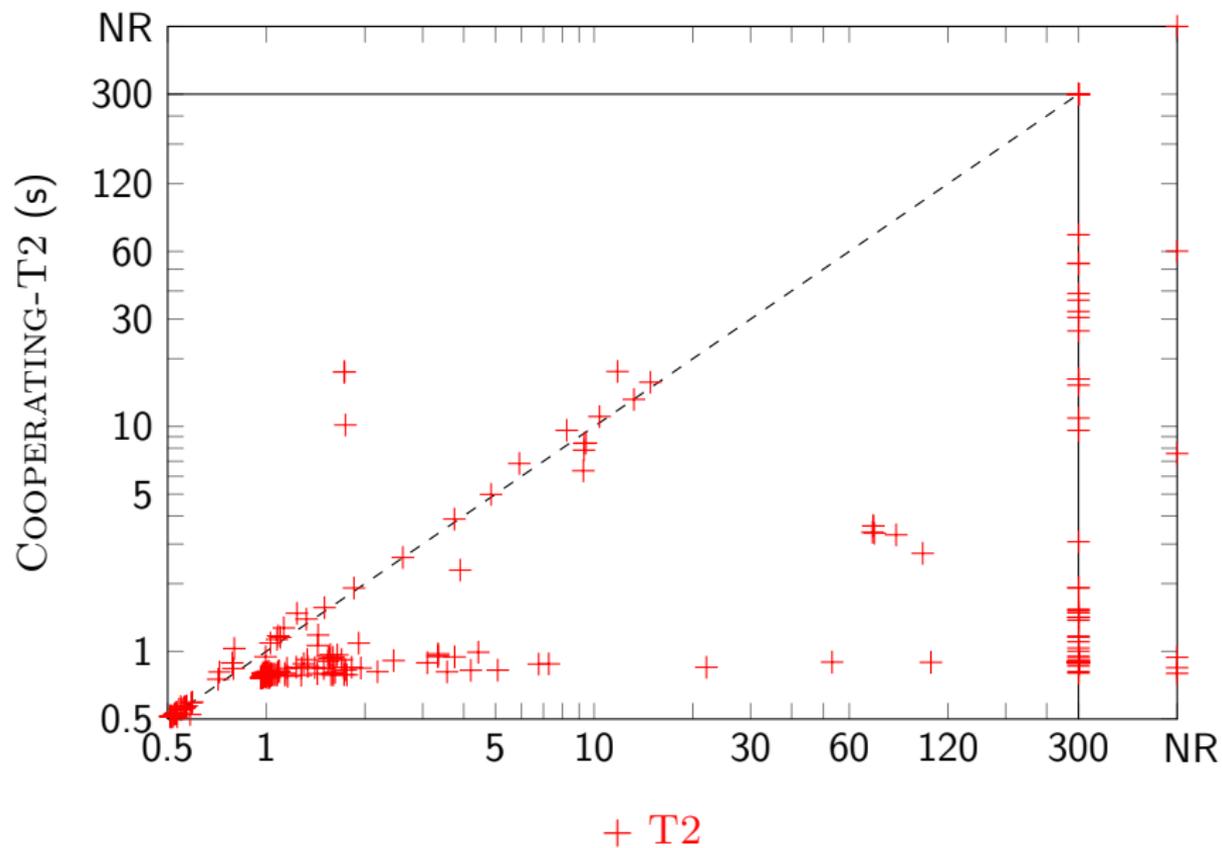
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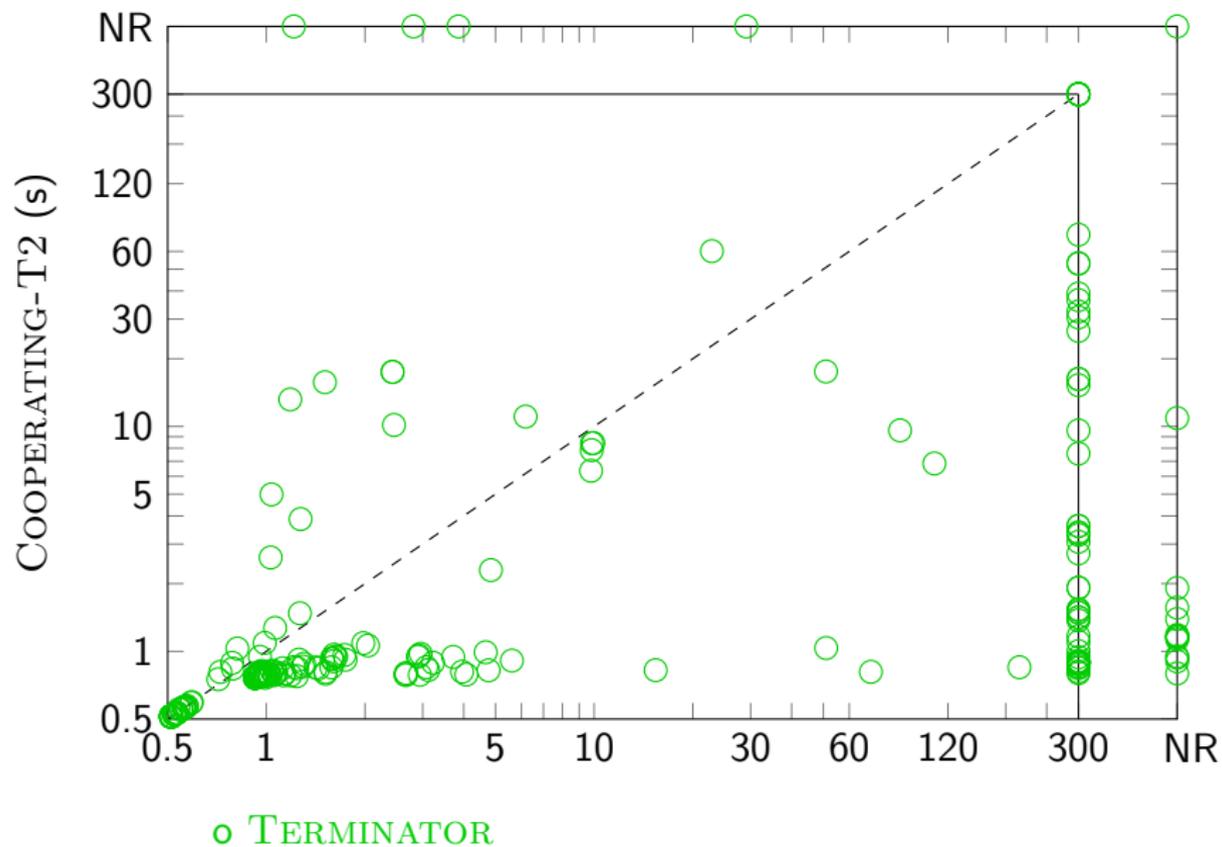
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	Term (#)	Term (avg. s)
COOPERATING-T2	245	3.42
APROVE	197	2.21
KITTEL	196	4.65
T2	189	5.15
APROVE+INTERPROC	185	1.53
TERMINATOR	177	4.99
SIZE-CHANGE/MCNP	156	17.50
ARMC	138	16.16

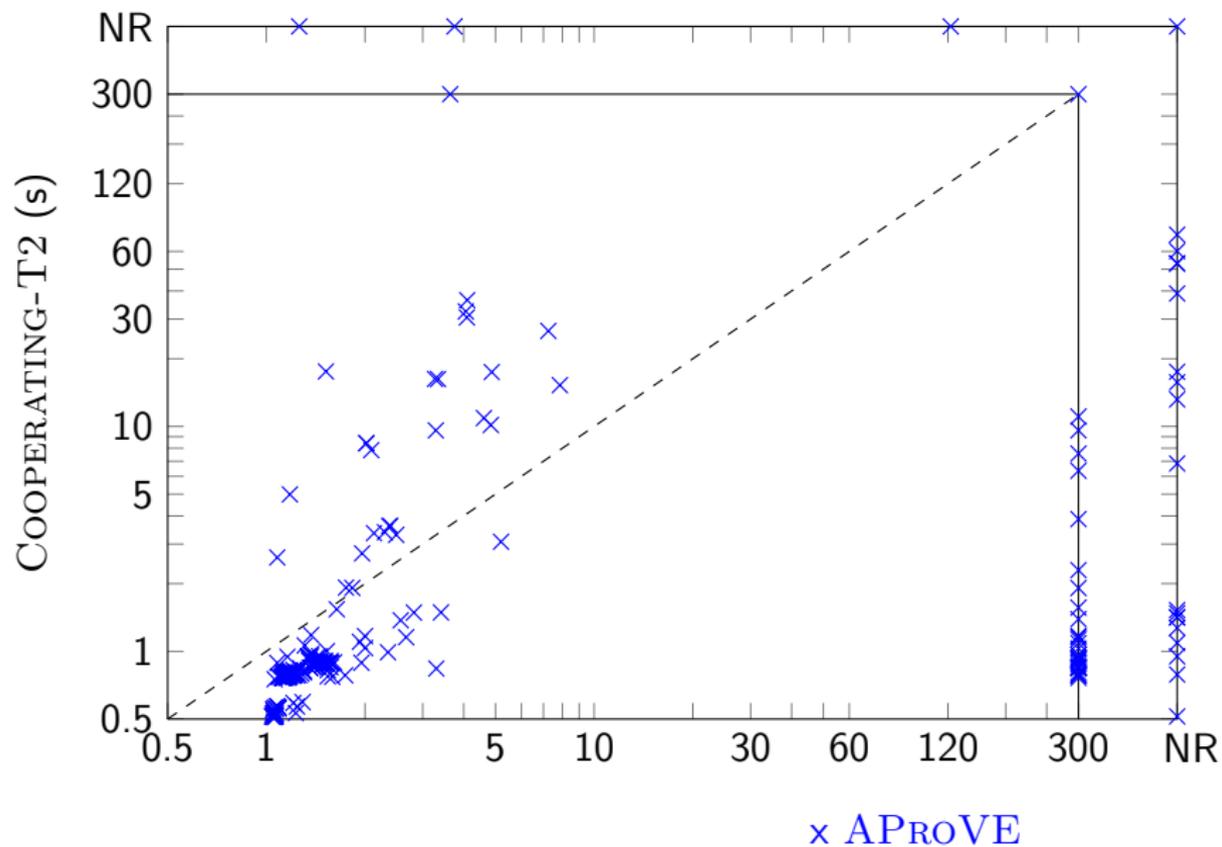
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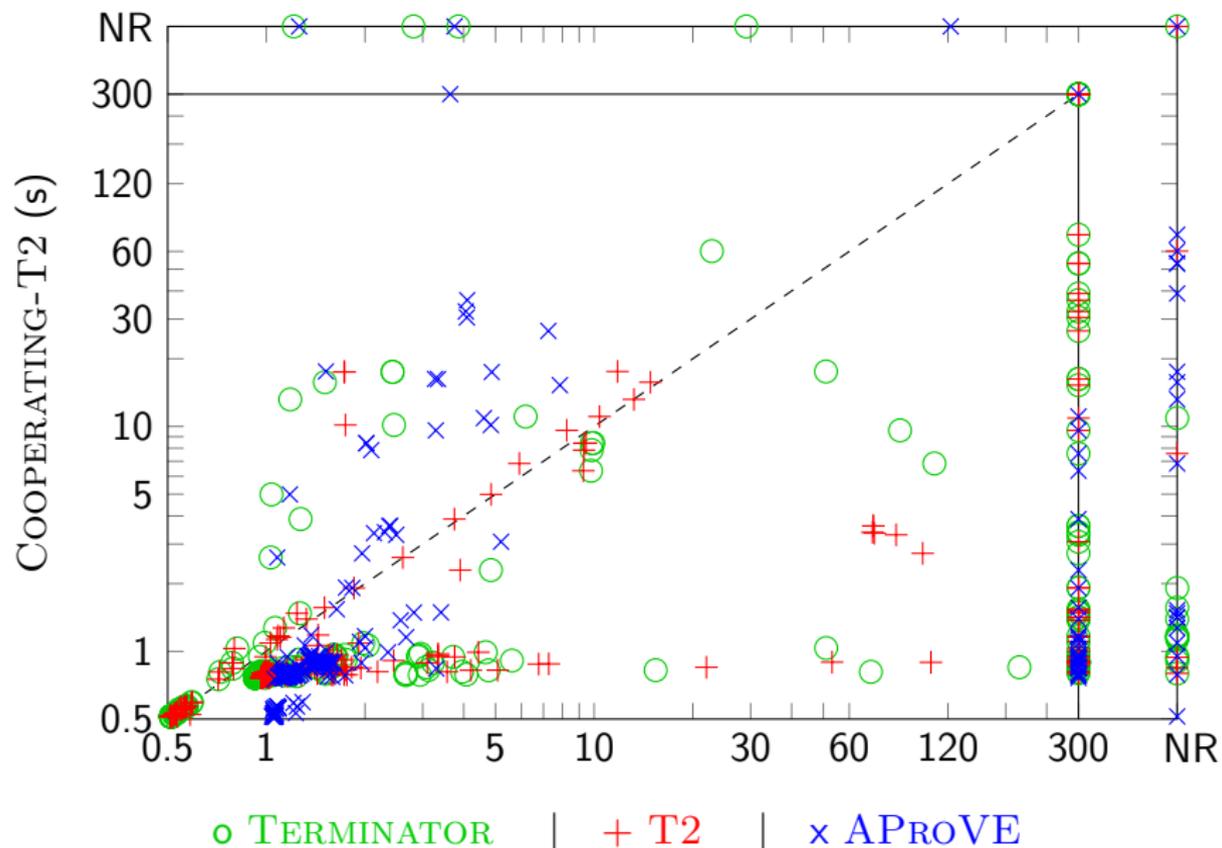
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Sources available: <http://research.microsoft.com/en-us/projects/t2/>