Semantics by lub's of Approximations

```
fact :: Int -> Int
fact  = \x -> if x <= 0 then 1 else fact(x-1) * x
```

Regard non-recursively defined approximations

$$\begin{aligned}
&\texttt{fact}_0 = \texttt{\textbackslash x -> bot}\\
&\texttt{fact}_1 = \texttt{\textbackslash x -> if x <= 0 then 1 else fact}_0\texttt{(x} - \texttt{1)} * \texttt{x}\\
&\texttt{fact}_2 = \texttt{\textbackslash x -> if x <= 0 then 1 else fact}_1\texttt{(x} - \texttt{1)} * \texttt{x}\\
&\qquad \ldots
\end{aligned}$$

Thus: $\texttt{fact}_{i+1} = \texttt{ff fact}_i$ where

```
ff :: (Int -> Int) -> (Int -> Int)
ff g  = \x -> if x <= 0 then 1 else g(x-1) * x
```

Define semantics fact of `fact` as

$$\mathit{fact} = \sqcup\{\mathit{fact}_0, \mathit{fact}_1, \ldots\} = \sqcup\{\mathit{ff}^i(\bot) \mid i \in \mathbb{N}\}$$

Semantics by Least Fixpoints

```
fact :: Int -> Int
fact  = \x -> if x <= 0 then 1 else fact(x-1) * x
```

Semantics $fact$ of `fact` should *satisfy* the defining equation

$$fact = \underbrace{\backslash\texttt{x -> if x <= 0 then 1 else}\ fact(\texttt{x} - 1) * \texttt{x}}_{ff\,(fact)}$$

where

```
ff :: (Int -> Int) -> (Int -> Int)
ff g  = \x -> if x <= 0 then 1 else g(x-1) * x
```

$fact$ should be undefined unless program enforces its definedness.
Define semantics $fact$ of `fact` as

$$fact = \mathrm{lfp}\ ff$$

Fixpoint Theorem

Both semantics are equivalent:

$$\mathrm{lfp}\, \mathit{ff} \quad = \quad \sqcup\{\mathit{ff}^i(\bot) \mid i \in \mathbb{N}\}$$

Thm. 2.1.17 (*Fixpoint Theorem*)
Let $\sqsubseteq$ be a cpo on D and let $f : D \to D$ be continuous.
Then f has a least fixpoint and we have $\mathrm{lfp}\, f \;=\; \sqcup\{f^i(\bot) | i \in \mathbb{N}\}$.

Proof.
$\bot \sqsubseteq f(\bot) \sqsubseteq f^2(\bot) \sqsubseteq f^3(\bot) \sqsubseteq \ldots$ by monotonicity of f

So $\{f^i(\bot) | i \in \mathbb{N}\}$ is a chain
and $\sqcup\{f^i(\bot) | i \in \mathbb{N}\}$ exists by completeness of $\sqsubseteq$.

To show: $\sqcup\{f^i(\bot) | i \in \mathbb{N}\}$ is the least fixpoint of f