

A Dependency Pair Framework for Innermost Complexity Analysis of Term Rewrite Systems

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joint work with [Lars Noschinski](#) and [Fabian Emmes](#)

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(cannot exploit strength of DPs for innermost rewriting)
 - **new approach: direct adaption of DP framework (CADE '11)**
 - modular combination of different techniques
 - automated and more powerful than previous approaches

Innermost Runtime Complexity

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- Example: $\iota_{\mathcal{R}} = \text{Pol}_1$

Dependency Tuples

$m(x, y) \rightarrow \text{if}(\text{gt}(x, y), x, y)$	$\text{gt}(0, k) \rightarrow \text{false}$	$p(0) \rightarrow 0$
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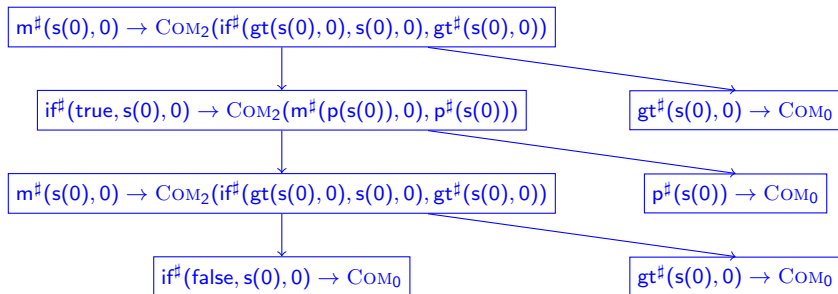
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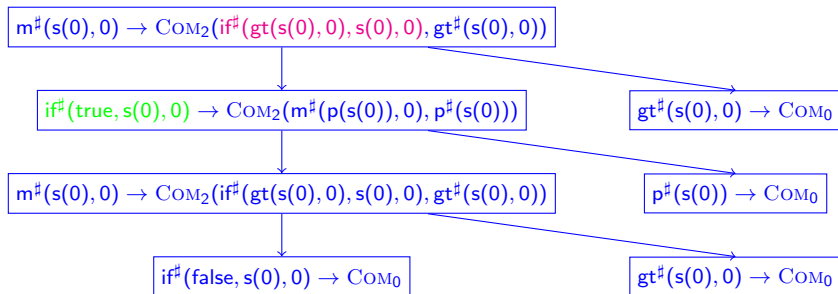


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$(\mathcal{D}, \mathcal{R})$ -Chain Tree:

Edge $\sigma_1(u^\# \rightarrow \text{COM}_n(v_1^\#, \dots, v_n^\#))$ to $\sigma_2(w^\# \rightarrow \text{COM}_m(\dots))$ if $v_i^\# \sigma_1 \xrightarrow{i}_{\mathcal{R}}^* w^\# \sigma_2$

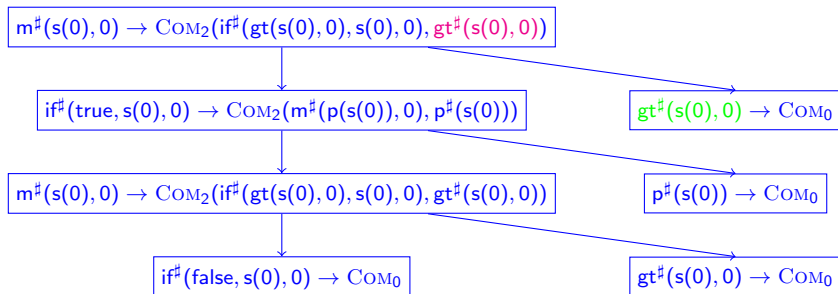


Chain Trees

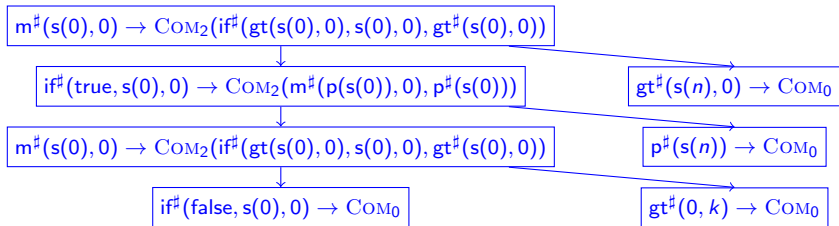
$$\begin{aligned} \text{DT}(\mathcal{R}) : \quad & m^\sharp(x, y) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(x, y), x, y), \text{gt}^\sharp(x, y)) & p^\sharp(0) & \rightarrow \text{COM}_0 \\ & \text{if}^\sharp(\text{true}, x, y) \rightarrow \text{COM}_2(m^\sharp(p(x), y), p^\sharp(x)) & p^\sharp(s(n)) & \rightarrow \text{COM}_0 \\ & \text{if}^\sharp(\text{false}, x, y) \rightarrow \text{COM}_0 & \text{gt}^\sharp(0, k) & \rightarrow \text{COM}_0 \\ & & \text{gt}^\sharp(s(n), 0) & \rightarrow \text{COM}_0 \\ & & \text{gt}^\sharp(s(n), s(k)) & \rightarrow \text{COM}_1(\text{gt}^\sharp(n, k)) \end{aligned}$$

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Edge $\sigma_1(u^\sharp \rightarrow \text{COM}_n(v_1^\sharp, \dots, v_n^\sharp))$ to $\sigma_2(w^\sharp \rightarrow \text{COM}_m(\dots))$ if $v_i^\sharp \sigma_1 \xrightarrow{i}^*_{\mathcal{R}} w^\sharp \sigma_2$

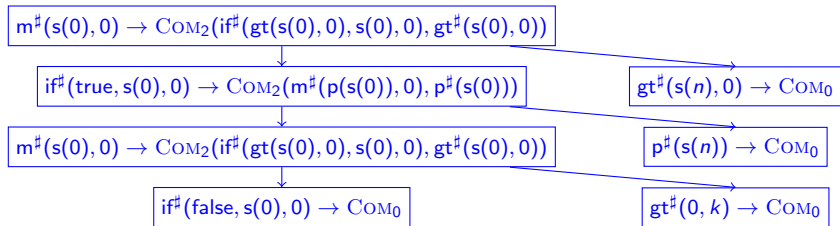


Chain Trees and Complexity



\mathcal{L}_R : length of longest $\xrightarrow{i} \mathcal{R}$ -sequence for $|t| \leq n$

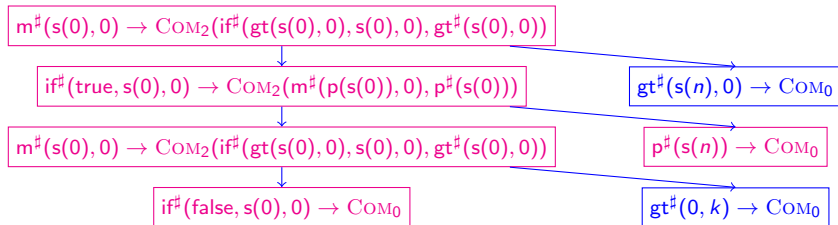
Chain Trees and Complexity



$\ell_{\mathcal{R}}$: length of longest $\xrightarrow{i}_{\mathcal{R}}$ -sequence for $|t| \leq n$

$\ell_{\langle \mathcal{D}, \mathcal{R} \rangle}$: maximal number of nodes
in chain tree with root $t^{\#} \rightarrow \text{COM}(\dots)$ for $|t| \leq n$

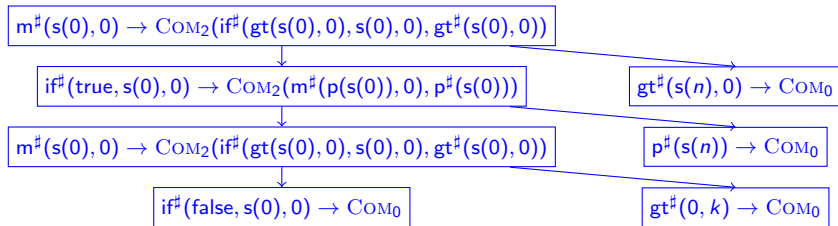
Chain Trees and Complexity



$\ell_{\mathcal{R}}$: length of longest $\xrightarrow{i}_{\mathcal{R}}$ -sequence for $|t| \leq n$

$\ell_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}$: maximal number of nodes from \mathcal{S}
in chain tree with root $t^\# \rightarrow \text{COM}(\dots)$ for $|t| \leq n$

Chain Trees and Complexity



$l_{\mathcal{R}}$: length of longest $\xrightarrow{i}_{\mathcal{R}}$ -sequence for $|t| \leq n$

$l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}$: maximal number of nodes from \mathcal{S}
in chain tree with root $t^{\#} \rightarrow \text{COM}(\dots)$ for $|t| \leq n$

Theorem

If $\mathcal{D} = \text{DT}(\mathcal{R})$, then $l_{\mathcal{R}} \leq l_{\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle}$.

Chain Trees and Complexity

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\Rightarrow Find out $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}$

Chain Trees and Complexity

Theorem

If $\mathcal{D} = DT(\mathcal{R})$, then $l_{\mathcal{R}} \leq l_{\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle}$.

⇒ Find out $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}$

⇒ Repeatedly replace **DT problem** $\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle$ by *simpler* $\langle \mathcal{D}', \mathcal{S}', \mathcal{R}' \rangle$, examine $l_{\langle \mathcal{D}', \mathcal{S}', \mathcal{R}' \rangle}$

Chain Trees and Complexity

Theorem

If $\mathcal{D} = DT(\mathcal{R})$, then $l_{\mathcal{R}} \leq l_{\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle}$.

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⇒ Start with **canonical DT problem** $\langle DT(\mathcal{R}), DT(\mathcal{R}), \mathcal{R} \rangle$

Chain Trees and Complexity

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If $\mathcal{D} = DT(\mathcal{R})$, then $l_{\mathcal{R}} \leq l_{\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle}$.

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⇒ Start with **canonical DT problem** $\langle DT(\mathcal{R}), DT(\mathcal{R}), \mathcal{R} \rangle$

DT Processor: $Proc(P) = (c, P')$ P, P' DT problems, $c \in \{Pol_0, Pol_1, \dots\}$

Chain Trees and Complexity

Theorem

If $\mathcal{D} = DT(\mathcal{R})$, then $l_{\mathcal{R}} \leq l_{\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle}$.

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$$P \overset{c}{\rightsquigarrow} P'$$

Chain Trees and Complexity

Theorem

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DT Processor: $Proc(P) = (c, P')$ P, P' DT problems, $c \in \{Pol_0, Pol_1, \dots\}$
where $l_P \leq \max(c, l_{P'})$

Proof Chain: $P_0 \overset{c_1}{\rightsquigarrow} P_1$

$P_0 = \langle DT(\mathcal{R}), DT(\mathcal{R}), \mathcal{R} \rangle$ canonical

Chain Trees and Complexity

Theorem

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where $l_P \leq \max(c, l_{P'})$

Proof Chain: $P_0 \overset{c_1}{\rightsquigarrow} P_1 \overset{c_2}{\rightsquigarrow} P_2$

$P_0 = \langle DT(\mathcal{R}), DT(\mathcal{R}), \mathcal{R} \rangle$ canonical

Chain Trees and Complexity

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Proof Chain: $P_0 \overset{c_1}{\rightsquigarrow} P_1 \overset{c_2}{\rightsquigarrow} P_2 \overset{c_3}{\rightsquigarrow} \dots \overset{c_k}{\rightsquigarrow} P_k$

$P_0 = \langle DT(\mathcal{R}), DT(\mathcal{R}), \mathcal{R} \rangle$ canonical

$P_k = \langle \mathcal{D}_k, \emptyset, \mathcal{R}_k \rangle$ solved

Chain Trees and Complexity

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 $l_{P_0} \leq \max(c_1, l_{P_1})$

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 $l_{P_0} \leq \max(c_1, c_2, l_{P_2})$

$P_0 = \langle DT(\mathcal{R}), DT(\mathcal{R}), \mathcal{R} \rangle$ canonical
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 $l_{\mathcal{R}} \leq l_{P_0} \leq \max(c_1, c_2, \dots, c_k)$ $P_k = \langle \mathcal{D}_k, \emptyset, \mathcal{R}_k \rangle$ solved

Leaf Removal Processor

Dependency Graph: edge from DT $u \rightarrow v$ to $w \rightarrow t$ in dep. graph

Leaf Removal Processor

Dependency Graph: edge from DT $u \rightarrow v$ to $w \rightarrow t$ in dep. graph iff
edge from $\sigma_1(u \rightarrow v)$ to $\sigma_2(w \rightarrow t)$ in chain tree

Leaf Removal Processor

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Example:

$\mathcal{R}: q(0, s(y), s(z)) \rightarrow 0, q(s(x), s(y), z) \rightarrow q(x, y, z), q(x, 0, s(z)) \rightarrow s(q(x, s(z), s(z)))$

Leaf Removal Processor

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$\mathcal{D}: q^\#(0, s(y), s(z)) \rightarrow \text{COM}_0 \qquad q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$

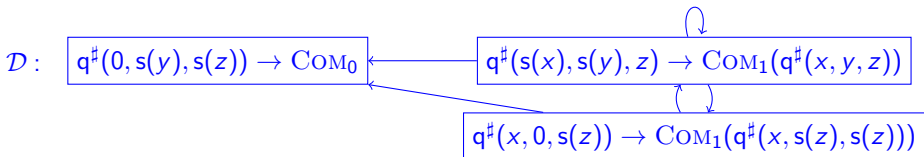
$q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$

Leaf Removal Processor

Dependency Graph: edge from DT $u \rightarrow v$ to $w \rightarrow t$ in dep. graph iff
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Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle$

\mathcal{R} : $q(0, s(y), s(z)) \rightarrow 0$, $q(s(x), s(y), z) \rightarrow q(x, y, z)$, $q(x, 0, s(z)) \rightarrow s(q(x, s(z), s(z)))$



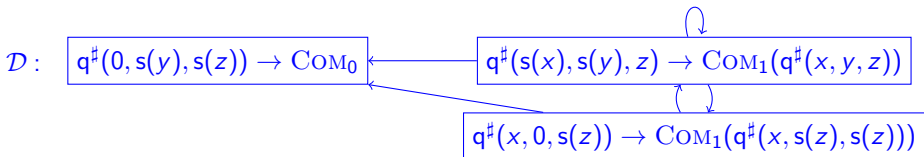
Leaf Removal Processor

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Leaf Removal Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D} \setminus \{w \rightarrow t\}, \mathcal{S} \setminus \{w \rightarrow t\}, \mathcal{R} \rangle$
if $w \rightarrow t$ is leaf in dependency graph

Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle$

\mathcal{R} : $q(0, s(y), s(z)) \rightarrow 0$, $q(s(x), s(y), z) \rightarrow q(x, y, z)$, $q(x, 0, s(z)) \rightarrow s(q(x, s(z), s(z)))$



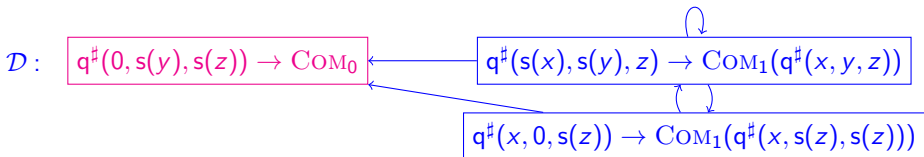
Leaf Removal Processor

Dependency Graph: edge from DT $u \rightarrow v$ to $w \rightarrow t$ in **dep. graph** iff
edge from $\sigma_1(u \rightarrow v)$ to $\sigma_2(w \rightarrow t)$ in **chain tree**

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\mathcal{R} : $q(0, s(y), s(z)) \rightarrow 0$, $q(s(x), s(y), z) \rightarrow q(x, y, z)$, $q(x, 0, s(z)) \rightarrow s(q(x, s(z), s(z)))$



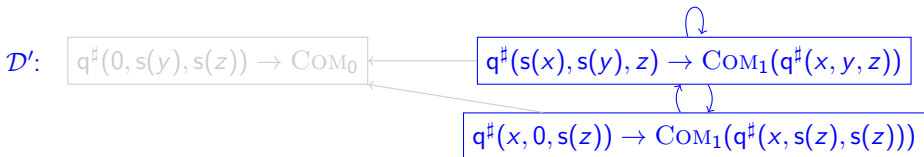
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Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{R} \rangle$

\mathcal{R} : $q(0, s(y), s(z)) \rightarrow 0$, $q(s(x), s(y), z) \rightarrow q(x, y, z)$, $q(x, 0, s(z)) \rightarrow s(q(x, s(z), s(z)))$



Usable Rules Processor

Usable Rules $\mathcal{U}_{\mathcal{R}}(\mathcal{D})$: rules from \mathcal{R} that can reduce rhs of \mathcal{D}

Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{R} \rangle$

\mathcal{R} : $q(0, s(y), s(z)) \rightarrow 0$, $q(s(x), s(y), z) \rightarrow q(x, y, z)$, $q(x, 0, s(z)) \rightarrow s(q(x, s(z), s(z)))$

\mathcal{D}' :

$q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$

$q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$

Usable Rules Processor

Usable Rules $\mathcal{U}_{\mathcal{R}}(\mathcal{D})$: rules from \mathcal{R} that can reduce rhs of \mathcal{D}

Usable Rules Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}, \mathcal{S}, \mathcal{U}_{\mathcal{R}}(\mathcal{D}) \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{R} \rangle$

\mathcal{R} : $q(0, s(y), s(z)) \rightarrow 0$, $q(s(x), s(y), z) \rightarrow q(x, y, z)$, $q(x, 0, s(z)) \rightarrow s(q(x, s(z), s(z)))$

\mathcal{D}' :

$q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$

$q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$

Usable Rules Processor

Usable Rules $\mathcal{U}_{\mathcal{R}}(\mathcal{D})$: rules from \mathcal{R} that can reduce rhs of \mathcal{D}

Usable Rules Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}, \mathcal{S}, \mathcal{U}_{\mathcal{R}}(\mathcal{D}) \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{U}_{\mathcal{R}}(\mathcal{D}') \rangle$

\mathcal{R} : $q(0, s(y), s(z)) \rightarrow 0$, $q(s(x), s(y), z) \rightarrow q(x, y, z)$, $q(x, 0, s(z)) \rightarrow s(q(x, s(z), s(z)))$

\mathcal{D}' :

$q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$

$q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$

Usable Rules Processor

Usable Rules $\mathcal{U}_{\mathcal{R}}(\mathcal{D})$: rules from \mathcal{R} that can reduce rhs of \mathcal{D}

Usable Rules Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}, \mathcal{S}, \mathcal{U}_{\mathcal{R}}(\mathcal{D}) \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{U}_{\mathcal{R}}(\mathcal{D}') \rangle$

$\mathcal{U}_{\mathcal{R}}(\mathcal{D}')$:

\mathcal{D}' :

$q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$

$q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$

Usable Rules Processor

Usable Rules $\mathcal{U}_{\mathcal{R}}(\mathcal{D})$: rules from \mathcal{R} that can reduce rhs of \mathcal{D}

Usable Rules Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}, \mathcal{S}, \mathcal{U}_{\mathcal{R}}(\mathcal{D}) \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset \rangle$

$\mathcal{U}_{\mathcal{R}}(\mathcal{D}')$:

\mathcal{D}' :

$q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$

$q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$

Extended DT Problems

DT Problem: $\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset \rangle$

\mathcal{D}' :

$q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$

$q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$

Extended DT Problems

DT Problem: $\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle$

Canonical

DT Problem: $\langle \text{DT}(\mathcal{R}), \text{DT}(\mathcal{R}), \mathcal{R} \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset \rangle$

\mathcal{D}' :

$$\text{q}^\#(s(x), s(y), z) \rightarrow \text{COM}_1(\text{q}^\#(x, y, z))$$

$$\text{q}^\#(x, 0, s(z)) \rightarrow \text{COM}_1(\text{q}^\#(x, s(z), s(z)))$$

Extended DT Problems

DT Problem: $\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle$

- when computing $\iota_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}$, we already took $\iota_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle}$ into account

Canonical

DT Problem: $\langle \text{DT}(\mathcal{R}), \text{DT}(\mathcal{R}), \mathcal{R} \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset \rangle$

\mathcal{D}' :

$$\text{q}^\#(s(x), s(y), z) \rightarrow \text{COM}_1(\text{q}^\#(x, y, z))$$

$$\text{q}^\#(x, 0, s(z)) \rightarrow \text{COM}_1(\text{q}^\#(x, s(z), s(z)))$$

Extended DT Problems

Extended DT Problem: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle$

- when computing $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}$, we already took $l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle}$ into account

Canonical

DT Problem: $\langle \text{DT}(\mathcal{R}), \text{DT}(\mathcal{R}), \mathcal{R} \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset \rangle$

\mathcal{D}' :

$$\boxed{q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))}$$

$$\boxed{q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))}$$

Extended DT Problems

Extended DT Problem: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle$

- when computing $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}$, we already took $l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle}$ into account
- $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle} = \begin{cases} l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}, & \text{if } l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle} > l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle} \end{cases}$

Canonical

DT Problem: $\langle \text{DT}(\mathcal{R}), \text{DT}(\mathcal{R}), \mathcal{R} \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle \xrightarrow{\text{Polo}} \langle \mathcal{D}', \mathcal{D}', \mathcal{R} \rangle \xrightarrow{\text{Polo}} \langle \mathcal{D}', \mathcal{D}', \emptyset \rangle$

\mathcal{D}' :

$$\text{q}^\#(s(x), s(y), z) \rightarrow \text{COM}_1(\text{q}^\#(x, y, z))$$

$$\text{q}^\#(x, 0, s(z)) \rightarrow \text{COM}_1(\text{q}^\#(x, s(z), s(z)))$$

Extended DT Problems

Extended DT Problem: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle$

- when computing $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}$, we already took $l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle}$ into account
- $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle} = \begin{cases} l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}, & \text{if } l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle} > l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle} \\ \text{Pol}_0, & \text{if } l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle} \leq l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle} \end{cases}$

Canonical

DT Problem: $\langle \text{DT}(\mathcal{R}), \text{DT}(\mathcal{R}), \mathcal{R} \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset \rangle$

\mathcal{D}' :

$$\text{q}^\#(s(x), s(y), z) \rightarrow \text{COM}_1(\text{q}^\#(x, y, z))$$

$$\text{q}^\#(x, 0, s(z)) \rightarrow \text{COM}_1(\text{q}^\#(x, s(z), s(z)))$$

Extended DT Problems

Extended DT Problem: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle$

- when computing $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}$, we already took $l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle}$ into account
- $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle} = \begin{cases} l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}, & \text{if } l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle} > l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle} \\ \text{Pol}_0, & \text{if } l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle} \leq l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle} \end{cases}$

Canonical Extended DT Problem: $\langle \text{DT}(\mathcal{R}), \text{DT}(\mathcal{R}), \emptyset, \mathcal{R} \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset \rangle$

\mathcal{D}' :

$$\text{q}^\#(s(x), s(y), z) \rightarrow \text{COM}_1(\text{q}^\#(x, y, z))$$

$$\text{q}^\#(x, 0, s(z)) \rightarrow \text{COM}_1(\text{q}^\#(x, s(z), s(z)))$$

Extended DT Problems

Extended DT Problem: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle$

- when computing $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}$, we already took $l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle}$ into account
- $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle} = \begin{cases} l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}, & \text{if } l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle} > l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle} \\ \text{Pol}_0, & \text{if } l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle} \leq l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle} \end{cases}$

Canonical Extended DT Problem: $\langle \text{DT}(\mathcal{R}), \text{DT}(\mathcal{R}), \emptyset, \mathcal{R} \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset \rangle$

\mathcal{D}' :

$$q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$$

$$q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$$

Extended DT Problems

Extended DT Problem: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle$

- when computing $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}$, we already took $l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle}$ into account
- $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle} = \begin{cases} l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}, & \text{if } l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle} > l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle} \\ \text{Pol}_0, & \text{if } l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle} \leq l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle} \end{cases}$

Canonical Extended DT Problem: $\langle \text{DT}(\mathcal{R}), \text{DT}(\mathcal{R}), \emptyset, \mathcal{R} \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset \rangle$

\mathcal{D}' :

$$\text{q}^\#(s(x), s(y), z) \rightarrow \text{COM}_1(\text{q}^\#(x, y, z))$$

$$\text{q}^\#(x, 0, s(z)) \rightarrow \text{COM}_1(\text{q}^\#(x, s(z), s(z)))$$

Extended DT Problems

Extended DT Problem: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle$

- when computing $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}$, we already took $l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle}$ into account
- $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle} = \begin{cases} l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}, & \text{if } l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle} > l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle} \\ \text{Pol}_0, & \text{if } l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle} \leq l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle} \end{cases}$

Canonical Extended DT Problem: $\langle \text{DT}(\mathcal{R}), \text{DT}(\mathcal{R}), \emptyset, \mathcal{R} \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$

\mathcal{D}' :

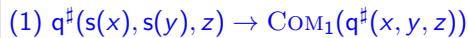
$$\text{q}^\#(s(x), s(y), z) \rightarrow \text{COM}_1(\text{q}^\#(x, y, z))$$

$$\text{q}^\#(x, 0, s(z)) \rightarrow \text{COM}_1(\text{q}^\#(x, s(z), s(z)))$$

Reduction Pair Processor

Termination: $\ell \succ r$ for all DPs and rules, remove DPs with $\ell \succ r$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$


$$(1) \text{q}^\#(s(x), s(y), z) \rightarrow \text{COM}_1(\text{q}^\#(x, y, z))$$


$$(2) \text{q}^\#(x, 0, s(z)) \rightarrow \text{COM}_1(\text{q}^\#(x, s(z), s(z)))$$

Reduction Pair Processor

Termination: $\ell \succsim r$ for all DPs and rules, remove DPs with $\ell \succ r$

Complexity: $\ell \succsim r$ for all DTs and rules, move DTs with $\ell \succ r$ from \mathcal{S} to \mathcal{K}

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$

(1) $q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$

(2) $q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$

Reduction Pair Processor

Termination: $\ell \succsim r$ for all DPs and rules, remove DPs with $\ell \succ r$

Complexity: $\ell \succsim r$ for all DTs and rules, move DTs with $\ell \succ r$ from \mathcal{S} to \mathcal{K}

Reduction Pair Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_m} \langle \mathcal{D}, \mathcal{S} \setminus \mathcal{D}_\succ, \mathcal{K} \cup \mathcal{D}_\succ, \mathcal{R} \rangle$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$

(1) $q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$

(2) $q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$

Reduction Pair Processor

Termination: $\ell \succ r$ for all DPs and rules, remove DPs with $\ell \succ r$

Complexity: $\ell \succ r$ for all DTs and rules, move DTs with $\ell \succ r$ from \mathcal{S} to \mathcal{K}

Reduction Pair Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_m} \langle \mathcal{D}, \mathcal{S} \setminus \mathcal{D}_\succ, \mathcal{K} \cup \mathcal{D}_\succ, \mathcal{R} \rangle$ if

- $\mathcal{D} \subseteq \succ \cup \succ$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$

(1) $q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$

(2) $q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$

Reduction Pair Processor

Termination: $l \succ r$ for all DPs and rules, remove DPs with $l \succ r$

Complexity: $l \succ r$ for all DTs and rules, move DTs with $l \succ r$ from \mathcal{S} to \mathcal{K}

Reduction Pair Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_m} \langle \mathcal{D}, \mathcal{S} \setminus \mathcal{D}_\succ, \mathcal{K} \cup \mathcal{D}_\succ, \mathcal{R} \rangle$ if

- $\mathcal{D} \subseteq \succ \cup \succ, \mathcal{R} \subseteq \succ$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$

(1) $q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$

(2) $q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$

Reduction Pair Processor

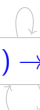
Termination: $l \succ r$ for all DPs and rules, remove DPs with $l \succ r$

Complexity: $l \succ r$ for all DTs and rules, move DTs with $l \succ r$ from \mathcal{S} to \mathcal{K}

Reduction Pair Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_m} \langle \mathcal{D}, \mathcal{S} \setminus \mathcal{D}_\succ, \mathcal{K} \cup \mathcal{D}_\succ, \mathcal{R} \rangle$ if

- $\mathcal{D} \subseteq \succ \cup \succ, \mathcal{R} \subseteq \succ$
- m is the maximal degree of polynomials $[f^\#]$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$


$$(1) \text{q}^\#(s(x), s(y), z) \rightarrow \text{COM}_1(\text{q}^\#(x, y, z))$$

$$(2) \text{q}^\#(x, 0, s(z)) \rightarrow \text{COM}_1(\text{q}^\#(x, s(z), s(z)))$$

Reduction Pair Processor

Termination: $\ell \succ r$ for all DPs and rules, remove DPs with $\ell \succ r$

Complexity: $\ell \succ r$ for all DTs and rules, move DTs with $\ell \succ r$ from \mathcal{S} to \mathcal{K}

Reduction Pair Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_m} \langle \mathcal{D}, \mathcal{S} \setminus \mathcal{D}_\succ, \mathcal{K} \cup \mathcal{D}_\succ, \mathcal{R} \rangle$ if

- $\mathcal{D} \subseteq \succ \cup \succ, \mathcal{R} \subseteq \succ$
- m is the maximal degree of polynomials $[f^\#]$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$

Polynomial Order

$$[\text{COM}_1](x) = x$$

$$[q^\#](x, y, z) = x$$

$$[s](x) = x + 1$$

$$(1) q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$$

$$(2) q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$$

Reduction Pair Processor

Termination: $\ell \succ r$ for all DPs and rules, remove DPs with $\ell \succ r$

Complexity: $\ell \succ r$ for all DTs and rules, move DTs with $\ell \succ r$ from \mathcal{S} to \mathcal{K}

Reduction Pair Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_m} \langle \mathcal{D}, \mathcal{S} \setminus \mathcal{D}_\succ, \mathcal{K} \cup \mathcal{D}_\succ, \mathcal{R} \rangle$ if

- $\mathcal{D} \subseteq \succ \cup \succ, \mathcal{R} \subseteq \succ$
- m is the maximal degree of polynomials $[f^\#]$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$

Polynomial Order

$$[\text{COM}_1](x) = x$$

$$[q^\#](x, y, z) = x$$

$$[s](x) = x + 1$$

$$(1) \quad q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$$

$$(2) \quad q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$$

Reduction Pair Processor

Termination: $\ell \succ r$ for all DPs and rules, remove DPs with $\ell \succ r$

Complexity: $\ell \succ r$ for all DTs and rules, move DTs with $\ell \succ r$ from \mathcal{S} to \mathcal{K}

Reduction Pair Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_m} \langle \mathcal{D}, \mathcal{S} \setminus \mathcal{D}_\succ, \mathcal{K} \cup \mathcal{D}_\succ, \mathcal{R} \rangle$ if

- $\mathcal{D} \subseteq \succ \cup \succ, \mathcal{R} \subseteq \succ$
- m is the maximal degree of polynomials $[f^\#]$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$

Polynomial Order

$$[\text{COM}_1](x) = x$$

$$[q^\#](x, y, z) = x$$

$$[s](x) = x + 1$$

$$(1) \quad q^\#(s(x), s(y), z) \succ \text{COM}_1(q^\#(x, y, z))$$

$$(2) \quad q^\#(x, 0, s(z)) \succ \text{COM}_1(q^\#(x, s(z), s(z)))$$

Reduction Pair Processor

Termination: $\ell \succ r$ for all DPs and rules, remove DPs with $\ell \succ r$

Complexity: $\ell \succ r$ for all DTs and rules, move DTs with $\ell \succ r$ from \mathcal{S} to \mathcal{K}

Reduction Pair Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_m} \langle \mathcal{D}, \mathcal{S} \setminus \mathcal{D}_\succ, \mathcal{K} \cup \mathcal{D}_\succ, \mathcal{R} \rangle$ if

- $\mathcal{D} \subseteq \succ \cup \succ$, $\mathcal{R} \subseteq \succ$
- m is the maximal degree of polynomials $[f^\#]$

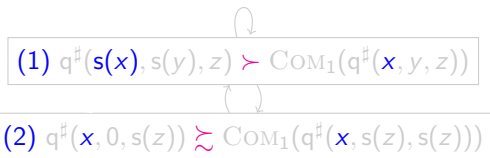
Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$
 $\xrightarrow{\text{Pol}_1} \langle \mathcal{D}', \{(2)\}, \{(1)\}, \emptyset \rangle$

Polynomial Order

$$[\text{COM}_1](x) = x$$

$$[q^\#](x, y, z) = x$$

$$[s](x) = x + 1$$



Knowledge Propagation Processor

Lemma: $l\langle \mathcal{D}, \{w \rightarrow t\}, \mathcal{R} \rangle \leq l\langle \mathcal{D}, \text{Pre}(w \rightarrow t), \mathcal{R} \rangle$

- $\text{Pre}(w \rightarrow t)$: all predecessors of $w \rightarrow t$ in dependency graph

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$
 $\xrightarrow{\text{Pol}_1} \langle \mathcal{D}', \{(2)\}, \{(1)\}, \emptyset \rangle$

(1) $q^\#(s(x), s(y), z) \rightarrow \text{COM}_1(q^\#(x, y, z))$

(2) $q^\#(x, 0, s(z)) \rightarrow \text{COM}_1(q^\#(x, s(z), s(z)))$

Knowledge Propagation Processor

Lemma: $l\langle \mathcal{D}, \{w \rightarrow t\}, \mathcal{R} \rangle \leq l\langle \mathcal{D}, \text{Pre}(w \rightarrow t), \mathcal{R} \rangle$

- $\text{Pre}(w \rightarrow t)$: all predecessors of $w \rightarrow t$ in dependency graph

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$
 $\xrightarrow{\text{Pol}_1} \langle \mathcal{D}', \{(2)\}, \{(1)\}, \emptyset \rangle$

$(1) \text{ q}^\#(s(x), s(y), z) \rightarrow \text{COM}_1(\text{q}^\#(x, y, z))$

$(2) \text{ q}^\#(x, 0, s(z)) \rightarrow \text{COM}_1(\text{q}^\#(x, s(z), s(z)))$

Knowledge Propagation Processor

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$\text{Pre}((2)) = \{(1)\}$

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Knowledge Propagation Processor

Lemma: $l_{\langle \mathcal{D}, \{w \rightarrow t\}, \mathcal{R} \rangle} \leq l_{\langle \mathcal{D}, \text{Pre}(w \rightarrow t), \mathcal{R} \rangle}$

- $\text{Pre}(w \rightarrow t)$: all predecessors of $w \rightarrow t$ in dependency graph
- $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle$: do not take $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle}$ into account if $l_{\langle \mathcal{D}, \mathcal{S}, \mathcal{R} \rangle} \leq l_{\langle \mathcal{D}, \mathcal{K}, \mathcal{R} \rangle}$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$
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KP Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}, \mathcal{S} \setminus \{w \rightarrow t\}, \mathcal{K}, \mathcal{R} \rangle$
 if $w \rightarrow t \in \mathcal{S}$ and $\text{Pre}(w \rightarrow t) \subseteq \mathcal{K}$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$
 $\xrightarrow{\text{Pol}_1} \langle \mathcal{D}', \{(2)\}, \{(1)\}, \emptyset \rangle$

$\text{Pre}((2)) = \{(1)\}$

Knowledge Propagation Processor

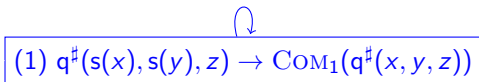
Lemma: $l_{\langle \mathcal{D}, \{w \rightarrow t\}, \mathcal{R} \rangle} \leq l_{\langle \mathcal{D}, \text{Pre}(w \rightarrow t), \mathcal{R} \rangle}$

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KP Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}, \mathcal{S} \setminus \{w \rightarrow t\}, \mathcal{K} \cup \{w \rightarrow t\}, \mathcal{R} \rangle$
 if $w \rightarrow t \in \mathcal{S}$ and $\text{Pre}(w \rightarrow t) \subseteq \mathcal{K}$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$
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$\text{Pre}((2)) = \{(1)\}$





Knowledge Propagation Processor

Lemma: $l_{\langle \mathcal{D}, \{w \rightarrow t\}, \mathcal{R} \rangle} \leq l_{\langle \mathcal{D}, \text{Pre}(w \rightarrow t), \mathcal{R} \rangle}$

- $\text{Pre}(w \rightarrow t)$: all predecessors of $w \rightarrow t$ in dependency graph
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KP Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}, \mathcal{S} \setminus \{w \rightarrow t\}, \mathcal{K} \cup \{w \rightarrow t\}, \mathcal{R} \rangle$
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 $\xrightarrow{\text{Pol}_1} \langle \mathcal{D}', \{(2)\}, \{(1)\}, \emptyset \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \emptyset, \{(1), (2)\}, \emptyset \rangle$

$\text{Pre}((2)) = \{(1)\}$

Knowledge Propagation Processor

Proof Chain: $P_0 \xrightarrow{c_1} \dots \xrightarrow{c_k} P_k$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$
 $\xrightarrow{Pol_1} \langle \mathcal{D}', \{(2)\}, \{(1)\}, \emptyset \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \emptyset, \{(1), (2)\}, \emptyset \rangle$

Knowledge Propagation Processor

Proof Chain: $P_0 \xrightarrow{c_1} \dots \xrightarrow{c_k} P_k$

$$l_{P_0} \leq \max(c_1, \dots, c_k)$$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$
 $\xrightarrow{Pol_1} \langle \mathcal{D}', \{(2)\}, \{(1)\}, \emptyset \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \emptyset, \{(1), (2)\}, \emptyset \rangle$

Knowledge Propagation Processor

Proof Chain: $P_0 \xrightarrow{c_1} \dots \xrightarrow{c_k} P_k$ $l_{\mathcal{R}} \leq l_{P_0} \leq \max(c_1, \dots, c_k)$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$
 $\xrightarrow{Pol_1} \langle \mathcal{D}', \{(2)\}, \{(1)\}, \emptyset \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \emptyset, \{(1), (2)\}, \emptyset \rangle$

Knowledge Propagation Processor

Proof Chain: $P_0 \xrightarrow{c_1} \dots \xrightarrow{c_k} P_k$ $l_{\mathcal{R}} \leq l_{P_0} \leq \max(c_1, \dots, c_k)$

\mathcal{R} : $q(0, s(y), s(z)) \rightarrow 0$, $q(s(x), s(y), z) \rightarrow q(x, y, z)$, $q(x, 0, s(z)) \rightarrow s(q(x, s(z), s(z)))$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$
 $\xrightarrow{Pol_1} \langle \mathcal{D}', \{(2)\}, \{(1)\}, \emptyset \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \emptyset, \{(1), (2)\}, \emptyset \rangle$

Knowledge Propagation Processor

Proof Chain: $P_0 \xrightarrow{c_1} \dots \xrightarrow{c_k} P_k$ $\iota_{\mathcal{R}} \leq \iota_{P_0} \leq \max(c_1, \dots, c_k)$

\mathcal{R} : $q(0, s(y), s(z)) \rightarrow 0$, $q(s(x), s(y), z) \rightarrow q(x, y, z)$, $q(x, 0, s(z)) \rightarrow s(q(x, s(z), s(z)))$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$
 $\xrightarrow{Pol_1} \langle \mathcal{D}', \{(2)\}, \{(1)\}, \emptyset \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \emptyset, \{(1), (2)\}, \emptyset \rangle$

$\iota_{\mathcal{R}} \leq \max(Pol_0, Pol_0, Pol_1, Pol_0)$

Knowledge Propagation Processor

Proof Chain: $P_0 \xrightarrow{c_1} \dots \xrightarrow{c_k} P_k$ $\iota_{\mathcal{R}} \leq \iota_{P_0} \leq \max(c_1, \dots, c_k)$

\mathcal{R} : $q(0, s(y), s(z)) \rightarrow 0$, $q(s(x), s(y), z) \rightarrow q(x, y, z)$, $q(x, 0, s(z)) \rightarrow s(q(x, s(z), s(z)))$

Example: $\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \mathcal{R} \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \mathcal{D}', \emptyset, \emptyset \rangle$
 $\xrightarrow{Pol_1} \langle \mathcal{D}', \{(2)\}, \{(1)\}, \emptyset \rangle \xrightarrow{Pol_0} \langle \mathcal{D}', \emptyset, \{(1), (2)\}, \emptyset \rangle$

$$\iota_{\mathcal{R}} \leq \max(Pol_0, Pol_0, Pol_1, Pol_0) = Pol_1$$

Narrowing Processor

Narrowing Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{S}', \mathcal{K}', \mathcal{R} \rangle$ where
in \mathcal{D}' , \mathcal{S}' , some $w \rightarrow t$ is replaced by all its narrowings

\mathcal{R} :

$m(x, y) \rightarrow \text{if}(\text{gt}(x, y), x, y)$	$\text{gt}(0, k) \rightarrow \text{false}$	$p(0) \rightarrow 0$
$\text{if}(\text{false}, x, y) \rightarrow 0$	$\text{gt}(s(n), 0) \rightarrow \text{true}$	$p(s(n)) \rightarrow n$
$\text{if}(\text{true}, x, y) \rightarrow s(m(p(x), y))$	$\text{gt}(s(n), s(k)) \rightarrow \text{gt}(n, k)$	

Narrowing Processor

Narrowing Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{S}', \mathcal{K}', \mathcal{R} \rangle$ where
in \mathcal{D}' , \mathcal{S}' , some $w \rightarrow t$ is replaced by all its narrowings

$\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle$

$\mathcal{R} :$

$m(x, y) \rightarrow$	$\text{if}(\text{gt}(x, y), x, y)$	$\text{gt}(0, k) \rightarrow$	false	$p(0) \rightarrow$	0
$\text{if}(\text{false}, x, y) \rightarrow$	0	$\text{gt}(s(n), 0) \rightarrow$	true	$p(s(n)) \rightarrow$	n
$\text{if}(\text{true}, x, y) \rightarrow$	$s(m(p(x), y))$	$\text{gt}(s(n), s(k)) \rightarrow$	$\text{gt}(n, k)$		

Narrowing Processor

Narrowing Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{S}', \mathcal{K}', \mathcal{R} \rangle$ where
in \mathcal{D}' , \mathcal{S}' , some $w \rightarrow t$ is replaced by all its narrowings

$$\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0^*} \langle \mathcal{D}_1, \mathcal{D}_1, \emptyset, \mathcal{R}_1 \rangle$$

$$\mathcal{R}_1 : \quad \begin{array}{ll} \text{gt}(0, k) \rightarrow \text{false} & \text{p}(0) \rightarrow 0 \\ \text{gt}(s(n), 0) \rightarrow \text{true} & \text{p}(s(n)) \rightarrow n \\ \text{gt}(s(n), s(k)) \rightarrow \text{gt}(n, k) & \end{array}$$

$$\mathcal{D}_1 : \quad m^\sharp(x, y) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(x, y), x, y), \text{gt}^\sharp(x, y))$$

$$\text{if}^\sharp(\text{true}, x, y) \rightarrow \text{COM}_2(m^\sharp(\text{p}(x), y), \text{p}^\sharp(x)) \quad \text{gt}^\sharp(s(n), s(k)) \rightarrow \text{COM}_1(\text{gt}^\sharp(n, k))$$

Narrowing Processor

Narrowing Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{S}', \mathcal{K}', \mathcal{R} \rangle$ where
in \mathcal{D}' , \mathcal{S}' , some $w \rightarrow t$ is replaced by all its narrowings

Narrowings of $m^\sharp(x, y) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(x, y), x, y), \text{gt}^\sharp(x, y))$

$\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0^*} \langle \mathcal{D}_1, \mathcal{D}_1, \emptyset, \mathcal{R}_1 \rangle$

\mathcal{R}_1 :

$\text{gt}(0, k) \rightarrow \text{false}$	$\text{p}(0) \rightarrow 0$
$\text{gt}(s(n), 0) \rightarrow \text{true}$	$\text{p}(s(n)) \rightarrow n$
$\text{gt}(s(n), s(k)) \rightarrow \text{gt}(n, k)$	

\mathcal{D}_1 : $m^\sharp(x, y) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(x, y), x, y), \text{gt}^\sharp(x, y))$

$\text{if}^\sharp(\text{true}, x, y) \rightarrow \text{COM}_2(m^\sharp(\text{p}(x), y), \text{p}^\sharp(x))$ $\text{gt}^\sharp(s(n), s(k)) \rightarrow \text{COM}_1(\text{gt}^\sharp(n, k))$

Narrowing Processor

Narrowing Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{S}', \mathcal{K}', \mathcal{R} \rangle$ where
in \mathcal{D}' , \mathcal{S}' , some $w \rightarrow t$ is replaced by all its narrowings

Narrowings of $m^\sharp(x, y) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(x, y), x, y), \text{gt}^\sharp(x, y))$

- $m^\sharp(0, k) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{false}, 0, k), \text{gt}^\sharp(0, k))$

$$\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0^*} \langle \mathcal{D}_1, \mathcal{D}_1, \emptyset, \mathcal{R}_1 \rangle$$

$$\mathcal{R}_1 : \quad \begin{array}{lll} \text{gt}(0, k) & \rightarrow & \text{false} & \text{p}(0) & \rightarrow & 0 \\ \text{gt}(s(n), 0) & \rightarrow & \text{true} & \text{p}(s(n)) & \rightarrow & n \\ \text{gt}(s(n), s(k)) & \rightarrow & \text{gt}(n, k) & & & \end{array}$$

$$\mathcal{D}_1 : \quad m^\sharp(x, y) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(x, y), x, y), \text{gt}^\sharp(x, y))$$

$$\text{if}^\sharp(\text{true}, x, y) \rightarrow \text{COM}_2(m^\sharp(\text{p}(x), y), \text{p}^\sharp(x)) \quad \text{gt}^\sharp(s(n), s(k)) \rightarrow \text{COM}_1(\text{gt}^\sharp(n, k))$$

Narrowing Processor

Narrowing Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{S}', \mathcal{K}', \mathcal{R} \rangle$ where
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Narrowings of $m^\sharp(x, y) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(x, y), x, y), \text{gt}^\sharp(x, y))$

- $m^\sharp(0, k) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{false}, 0, k), \text{gt}^\sharp(0, k))$
- $m^\sharp(s(n), 0) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{true}, s(n), 0), \text{gt}^\sharp(s(n), 0))$

$\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0^*} \langle \mathcal{D}_1, \mathcal{D}_1, \emptyset, \mathcal{R}_1 \rangle$

$\mathcal{R}_1 :$

$\text{gt}(0, k) \rightarrow \text{false}$	$p(0) \rightarrow 0$
$\text{gt}(s(n), 0) \rightarrow \text{true}$	$p(s(n)) \rightarrow n$
$\text{gt}(s(n), s(k)) \rightarrow \text{gt}(n, k)$	

$\mathcal{D}_1 : m^\sharp(x, y) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(x, y), x, y), \text{gt}^\sharp(x, y))$

$\text{if}^\sharp(\text{true}, x, y) \rightarrow \text{COM}_2(m^\sharp(p(x), y), p^\sharp(x))$ $\text{gt}^\sharp(s(n), s(k)) \rightarrow \text{COM}_1(\text{gt}^\sharp(n, k))$

Narrowing Processor

Narrowing Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{S}', \mathcal{K}', \mathcal{R} \rangle$ where
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Narrowings of $m^\sharp(x, y) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(x, y), x, y), \text{gt}^\sharp(x, y))$

- $m^\sharp(0, k) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{false}, 0, k), \text{gt}^\sharp(0, k))$
- $m^\sharp(s(n), 0) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{true}, s(n), 0), \text{gt}^\sharp(s(n), 0))$
- $m^\sharp(s(n), s(k)) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(n, k), s(n), s(k)), \text{gt}^\sharp(s(n), s(k)))$

$\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0^*} \langle \mathcal{D}_1, \mathcal{D}_1, \emptyset, \mathcal{R}_1 \rangle$

$\mathcal{R}_1 :$

$\text{gt}(0, k) \rightarrow \text{false}$	$p(0) \rightarrow 0$
$\text{gt}(s(n), 0) \rightarrow \text{true}$	$p(s(n)) \rightarrow n$
$\text{gt}(s(n), s(k)) \rightarrow \text{gt}(n, k)$	

$\mathcal{D}_1: m^\sharp(x, y) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(x, y), x, y), \text{gt}^\sharp(x, y))$

$\text{if}^\sharp(\text{true}, x, y) \rightarrow \text{COM}_2(m^\sharp(p(x), y), p^\sharp(x))$ $\text{gt}^\sharp(s(n), s(k)) \rightarrow \text{COM}_1(\text{gt}^\sharp(n, k))$

Narrowing Processor

Narrowing Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{S}', \mathcal{K}', \mathcal{R} \rangle$ where
in \mathcal{D}' , \mathcal{S}' , some $w \rightarrow t$ is replaced by all its narrowings

Narrowings of $m^\sharp(x, y) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(x, y), x, y), \text{gt}^\sharp(x, y))$

- $m^\sharp(s(n), 0) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{true}, s(n), 0), \text{gt}^\sharp(s(n), 0))$
- $m^\sharp(s(n), s(k)) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(n, k), s(n), s(k)), \text{gt}^\sharp(s(n), s(k)))$

$\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0^*} \langle \mathcal{D}_1, \mathcal{D}_1, \emptyset, \mathcal{R}_1 \rangle \xrightarrow{\text{Pol}_0^*} \langle \mathcal{D}_2, \mathcal{D}_2, \emptyset, \mathcal{R}_1 \rangle$

$\mathcal{R}_1 :$

$\text{gt}(0, k) \rightarrow \text{false}$	$\text{p}(0) \rightarrow 0$
$\text{gt}(s(n), 0) \rightarrow \text{true}$	$\text{p}(s(n)) \rightarrow n$
$\text{gt}(s(n), s(k)) \rightarrow \text{gt}(n, k)$	

$\mathcal{D}_2: m^\sharp(s(n), 0) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{true}, s(n), 0), \text{gt}^\sharp(s(n), 0))$

$m^\sharp(s(n), s(k)) \rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(n, k), s(n), s(k)), \text{gt}^\sharp(s(n), s(k)))$

$\text{if}^\sharp(\text{true}, x, y) \rightarrow \text{COM}_2(m^\sharp(\text{p}(x), y), \text{p}^\sharp(x))$ $\text{gt}^\sharp(s(n), s(k)) \rightarrow \text{COM}_1(\text{gt}^\sharp(n, k))$

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Narrowing Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0} \langle \mathcal{D}', \mathcal{S}', \mathcal{K}', \mathcal{R} \rangle$ where
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$$\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0^*} \langle \mathcal{D}_1, \mathcal{D}_1, \emptyset, \mathcal{R}_1 \rangle \xrightarrow{\text{Pol}_0^*} \langle \mathcal{D}_2, \mathcal{D}_2, \emptyset, \mathcal{R}_1 \rangle \xrightarrow{\text{Pol}_0^*} \langle \mathcal{D}_3, \mathcal{D}_3, \emptyset, \mathcal{R}_2 \rangle$$

$$\mathcal{R}_2 : \quad \begin{aligned} \text{gt}(0, k) &\rightarrow \text{false} \\ \text{gt}(s(n), 0) &\rightarrow \text{true} \\ \text{gt}(s(n), s(k)) &\rightarrow \text{gt}(n, k) \end{aligned}$$

$$\begin{aligned} \mathcal{D}_3 : \quad m^\sharp(s(n), 0) &\rightarrow \text{COM}_2(\text{if}^\sharp(\text{true}, s(n), 0), \text{gt}^\sharp(s(n), 0)) \\ m^\sharp(s(n), s(k)) &\rightarrow \text{COM}_2(\text{if}^\sharp(\text{gt}(n, k), s(n), s(k)), \text{gt}^\sharp(s(n), s(k))) \\ \text{if}^\sharp(\text{true}, s(n), y) &\rightarrow \text{COM}_2(m^\sharp(n, y), p^\sharp(s(n))) \quad \text{gt}^\sharp(s(n), s(k)) \rightarrow \text{COM}_1(\text{gt}^\sharp(n, k)) \end{aligned}$$

Narrowing Processor

Reduction Pair Processor: $\langle \mathcal{D}, \mathcal{S}, \mathcal{K}, \mathcal{R} \rangle \xrightarrow{\text{Pol}_m} \langle \mathcal{D}, \mathcal{S} \setminus \mathcal{D}_\succ, \mathcal{K} \cup \mathcal{D}_\succ, \mathcal{R} \rangle$ where m is the maximal degree of polynomials $[f^\#]$

Polynomial Order

- $[0] = [\text{true}] = [\text{false}] = [p^\#](x) = 0, \quad [s](x) = x + 2$
- $[\text{gt}](x, y) = [\text{gt}^\#](x, y) = x$
- $[m^\#](x, y) = (x + 1)^2, \quad [\text{if}^\#](x, y, z) = y^2$

$\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow{\text{Pol}_0^*} \langle \mathcal{D}_1, \mathcal{D}_1, \emptyset, \mathcal{R}_1 \rangle \xrightarrow{\text{Pol}_0^*} \langle \mathcal{D}_2, \mathcal{D}_2, \emptyset, \mathcal{R}_2 \rangle \xrightarrow{\text{Pol}_0^*} \langle \mathcal{D}_3, \mathcal{D}_3, \emptyset, \mathcal{R}_2 \rangle$

$\mathcal{R}_2 :$

$\text{gt}(0, k)$	\succ	false
$\text{gt}(s(n), 0)$	\succ	true
$\text{gt}(s(n), s(k))$	\succ	$\text{gt}(n, k)$

$\mathcal{D}_3 :$

- $m^\#(s(n), 0) \succ \text{COM}_2(\text{if}^\#(\text{true}, s(n), 0), \text{gt}^\#(s(n), 0))$
- $m^\#(s(n), s(k)) \succ \text{COM}_2(\text{if}^\#(\text{gt}(n, k), s(n), s(k)), \text{gt}^\#(s(n), s(k)))$
- $\text{if}^\#(\text{true}, s(n), y) \succ \text{COM}_2(m^\#(s(n), y), p^\#(s(n))) \quad \text{gt}^\#(s(n), s(k)) \succ \text{COM}_1(\text{gt}^\#(n, k))$

Narrowing Processor

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$$\xrightarrow{\text{Pol}_2} \langle \mathcal{D}_3, \emptyset, \mathcal{D}_3, \mathcal{R}_2 \rangle$$

$$\mathcal{R}_2 : \quad \begin{array}{l} \text{gt}(0, k) \rightsquigarrow \text{false} \\ \text{gt}(s(n), 0) \rightsquigarrow \text{true} \\ \text{gt}(s(n), s(k)) \rightsquigarrow \text{gt}(n, k) \end{array}$$

$$\begin{array}{l} \mathcal{D}_3 : \quad m^\#(s(n), 0) \succ \text{COM}_2(\text{if}^\#(\text{true}, s(n), 0), \text{gt}^\#(s(n), 0)) \\ \quad m^\#(s(n), s(k)) \succ \text{COM}_2(\text{if}^\#(\text{gt}(n, k), s(n), s(k)), \text{gt}^\#(s(n), s(k))) \\ \quad \text{if}^\#(\text{true}, s(n), y) \succ \text{COM}_2(m^\#(s(n), y), p^\#(s(n))) \quad \text{gt}^\#(s(n), s(k)) \succ \text{COM}_1(\text{gt}^\#(n, k)) \end{array}$$

Narrowing Processor

$$\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \xrightarrow[\sim]{Pol_0^*} \langle \mathcal{D}_1, \mathcal{D}_1, \emptyset, \mathcal{R}_1 \rangle \xrightarrow[\sim]{Pol_0^*} \langle \mathcal{D}_2, \mathcal{D}_2, \emptyset, \mathcal{R}_1 \rangle \xrightarrow[\sim]{Pol_0^*} \langle \mathcal{D}_3, \mathcal{D}_3, \emptyset, \mathcal{R}_2 \rangle$$
$$\xrightarrow[\sim]{Pol_2} \langle \mathcal{D}_3, \emptyset, \mathcal{D}_3, \mathcal{R}_2 \rangle$$

$$l_{\mathcal{R}} \leq \max(Pol_0, \dots, Pol_0, Pol_2)$$

Narrowing Processor

$$\langle \mathcal{D}, \mathcal{D}, \emptyset, \mathcal{R} \rangle \begin{array}{l} \overset{Pol_0^*}{\rightsquigarrow} \\ \overset{Pol_2}{\rightsquigarrow} \end{array} \langle \mathcal{D}_1, \mathcal{D}_1, \emptyset, \mathcal{R}_1 \rangle \begin{array}{l} \overset{Pol_0^*}{\rightsquigarrow} \\ \rightsquigarrow \end{array} \langle \mathcal{D}_2, \mathcal{D}_2, \emptyset, \mathcal{R}_1 \rangle \overset{Pol_0^*}{\rightsquigarrow} \langle \mathcal{D}_3, \mathcal{D}_3, \emptyset, \mathcal{R}_2 \rangle$$

$$l_{\mathcal{R}} \leq \max(Pol_0, \dots, Pol_0, Pol_2) = Pol_2$$

DT Framework for Innermost Complexity Analysis

- *Direct* adaption of DP framework for termination analysis

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- Experiments on 1323 TRSs from *Termination Problem Data Base*
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		CaT					
		Pol_0	Pol_1	Pol_2	Pol_3	no result	Σ
AProVE	Pol_0	-	182	-	-	27	209
	Pol_1	-	187	7	-	76	270
	Pol_2	-	32	2	-	83	117
	Pol_3	-	6	-	-	16	22
	no result	-	27	3	1	674	705
	Σ	0	434	12	1	876	1323

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		TCT					Σ
		Pol_0	Pol_1	Pol_2	Pol_3	no result	
AProVE	Pol_0	10	157	-	-	42	209
	Pol_1	-	152	1	-	117	270
	Pol_2	-	35	-	-	82	117
	Pol_3	-	5	-	-	17	22
	no result	-	22	3	-	680	705
	Σ	10	371	4	0	938	1323

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