

Ground Resolution Algorithm

Goal: Determine whether $\{\varphi_1, \dots, \varphi_k\} \models \varphi$ holds

1. Let ξ be the formula $\varphi_1 \wedge \dots \wedge \varphi_k \wedge \neg\varphi$.
2. Transform ξ into Skolem normal form $\forall X_1, \dots, X_n \psi$.
3. Transform ψ into CNF resp. into clause set $\mathcal{K}(\psi)$.
4. Choose an enumeration $\{K_1, K_2, \dots\}$ of all ground instances of the clauses from $\mathcal{K}(\psi)$.
5. Compute $Res^*(\{K_1\})$, $Res^*(\{K_1, K_2\})$, $Res^*(\{K_1, K_2, K_3\})$, \dots
If one of these sets contains \square , stop and return **“true”**.

Resolution for Predicate Logic

$$\{ \{p(X), \neg q(X)\}, \{\neg p(f(Y))\}, \{q(f(a))\} \}$$

- use substitution $\{X/f(Y)\}$ for resolution of the first two clauses
- $p(X)[X/f(Y)] = p(f(Y))$ and $\neg p(f(Y))[X/f(Y)] = \neg p(f(Y))$
- $\{X/f(Y)\}$ is *most general unifier* of $\{p(X), p(f(Y))\}$
- resolvent is $\{\neg q(X)[X/f(Y)]\} = \{\neg q(f(Y))\}$