

## Notes:

- Please solve these exercises in **groups of three or four!**
- The solutions must be handed in **directly before (very latest: at the beginning of)** the exercise course on Friday, June 23rd, 2017, in lecture hall **AH 2**. Alternatively you can drop your solutions into a box which is located right next to Prof. Giesl's office (this box is emptied **15 minutes before** the exercise course starts).
- Please write the **names** and **immatriculation numbers** of all students on your solution. Also please staple the individual sheets!

**Exercise 1 (Semantics):**
**(3 + 3 + 3 + 1 = 10 points)**

 Consider the following logic program  $\mathcal{P}$  over the signature  $(\Sigma, \Delta)$  with  $0, s \in \Sigma$  and  $\text{minus}, \text{pred} \in \Delta$ :

```

minus(X, 0, X).
minus(X, Y, Z) :- minus(A, B, Z), pred(X, A), pred(Y, B).
pred(s(s(X)), s(Y)) :- pred(s(X), Y).
pred(s(X), X).
    
```

Also consider the following query:

```
?- minus(s(s(s(0))), Y, s(0)).
```

## Hints:

- In each of the computation steps that you give in this exercise, you can abbreviate those parts of the substitution that did not change by „...“.
- Show a successful computation for the query above (i.e., a computation of the form  $(G, \emptyset) \vdash_{\mathcal{P}}^+ (\square, \sigma)$  where  $G = \{\neg \text{minus}(s^3(0), Y, s(0))\}$ ). Also give the answer substitution.
  - Show a finite unsuccessful computation for the query above (i.e., a computation of the form  $(G, \emptyset) \vdash_{\mathcal{P}}^+ (G_1, \sigma_1)$  where  $G = \{\neg \text{minus}(s^3(0), Y, s(0))\}$ ,  $G_1 \neq \square$  and there are no  $G_2$  and  $\sigma_2$  such that  $(G_1, \sigma_1) \vdash_{\mathcal{P}} (G_2, \sigma_2)$ ).
  - Indicate an infinite computation for the query above by giving the first few steps. Give enough steps so that it is obvious how the infinite computation looks like.
  - What is  $D[\mathcal{P}, G]$  in this example?

*Example:* The query

```
?- pred(s(s(0)), Y).
```

 has two successful derivations. Here we use variable renamings to replace the variables  $X$  and  $Y$  in the program clauses by  $X_1$  and  $Y_1$  resp.  $X_2$ :

$$\begin{aligned}
 (\{\underline{\neg \text{pred}(s(s(0)), Y)}\}, \emptyset) \vdash_{\mathcal{P}} (\{\underline{\neg \text{pred}(s(0), Y_1)}\}, \{X_1/0, Y/s(Y_1)\}) \\
 \vdash_{\mathcal{P}} (\square, \{X_1/0, Y/s(0), X_2/0, Y_1/0\})
 \end{aligned} \tag{1}$$

$$(\{\underline{\neg \text{pred}(s(s(0)), Y)}\}, \emptyset) \vdash_{\mathcal{P}} (\square, \{X_1/s(0), Y/s(0)\}) \tag{2}$$

 The answer substitution for both derivations is  $\{Y/s(0)\}$ .

**Exercise 2 (Fixpoints):**

**(3 + 3 + 3 = 9 points)**

Consider the function  $f: Pot(\mathbb{N}) \rightarrow Pot(\mathbb{N})$ .

$$f(M) = \begin{cases} \{0\}, & \text{if } M = \emptyset \\ \{(\sum_{y \in X} y) - \min X \mid \emptyset \neq X \subseteq M\}, & \text{if } M \neq \emptyset \text{ is finite} \\ \mathbb{N}, & \text{otherwise} \end{cases}$$

So for example,  $f(\{2, 5\}) = \{(\sum_{y \in \{2\}} y) - \min \{2\}, (\sum_{y \in \{5\}} y) - \min \{5\}, (\sum_{y \in \{2, 5\}} y) - \min \{2, 5\}\} = \{0, 5\}$ .

- a) Prove that  $f$  is monotonic.
- b) Prove or disprove that  $f$  is continuous.
- c) Please give all fixpoints of  $f$  and mark the least fixpoint.

**Exercise 3 (Fixpoint Semantics):**

**(3 + 1 + 1 + 1 = 6 points)**

Reconsider the logic program  $\mathcal{P}$  over the signature  $(\Sigma, \Delta)$  with  $0, s \in \Sigma$  and  $\text{minus}, \text{pred} \in \Delta$  from Exercise 1, where we remove the recursive rule of the predicate  $\text{pred}$ :

```
minus(X, 0, X).
minus(X, Y, Z) :- minus(A, B, Z), pred(X, A), pred(Y, B).
pred(s(X), X).
```

- a) For each  $i \in \mathbb{N}$  explicitly give  $\text{trans}_{\mathcal{P}}^i(\emptyset)$ .
- b) Compute the set  $\text{lfp}(\text{trans}_{\mathcal{P}})$ .
- c) Give  $F[\mathcal{P}, \{\neg \text{minus}(s(s(0)), Y, Z)\}]$ .
- d) Give  $F[\mathcal{P}, \{\neg \text{minus}(X, s(Y), Y)\}]$ .

**Exercise 4 (Proofs):**

**(2 + 2 = 4 points)**

- a) Please show that every continuous function  $f: Pot(M) \rightarrow Pot(M)$  is monotonic.
- b) Please show that for every *finite* chain

$$M_1 \subseteq M_2 \subseteq \dots \subseteq M_n$$

and every monotonic function  $f: Pot(M) \rightarrow Pot(M)$  we have:

$$f\left(\bigcup_{i=1}^n M_i\right) = \bigcup_{i=1}^n f(M_i)$$