

Notes:

- Please solve these exercises in **groups of two!**
- Please register at <https://aprove.informatik.rwth-aachen.de/tes11/> (https, not http!).
- The solutions must be handed in **directly before (very latest: at the beginning of)** the exercise course on Wednesday, April 20th, 2011, in lecture hall **AH 2**. Alternatively you can drop your solutions into a box which is located right next to Prof. Giesl's office (until the exercise course starts).
- Please write the **names** and **immatriculation numbers** of all (two) students on your solution. Please staple the individual sheets!

Exercise 1 (Syntax and Semantics):
(2 + 4 = 6 points)

- a) Give a set of equalities that describes the function of the arithmetic operators $-$ as minus and $/$ as div for subtraction and division on natural numbers. The exact semantics should be as follows:

$$\begin{aligned} \text{div}(x, y) &= \lceil x/y \rceil && \text{if } y > 0 \\ \text{minus}(x, y) &= \begin{cases} x - y & \text{if } x > y \\ 0 & \text{otherwise} \end{cases} \end{aligned}$$

Thus, $\text{minus}(15, 10) = 5$, $\text{minus}(5, 10) = 0$ and $\text{div}(15, 10) = 2$. You should not define div for a divisor of 0.

Use the representation of natural numbers presented in the lecture, where 0 is represented by $\mathcal{O} \in \Sigma_0$ and n is represented by applying the successor symbol $s \in \Sigma_1$ n times (i.e., by $s^n(\mathcal{O})$). Define the signature of your equalities explicitly by giving definitions for Σ_i and Σ .

- b) Let $\Sigma = \Sigma_0 \cup \Sigma_1 \cup \Sigma_2$ with $\Sigma_0 = \{\mathcal{O}\}$, $\Sigma_1 = \{s\}$ and $\Sigma_2 = \{\text{plus}\}$. Consider $\mathcal{E} = \{\text{plus}(\mathcal{O}, y) \equiv y, \text{plus}(s(x), y) \equiv s(\text{plus}(x, y))\}$, the set of equations describing $+$ on our representation of natural numbers. Prove that $\mathcal{E} \not\models \text{plus}(x, y) \equiv \text{plus}(y, x)$.

Hints:

- You can use a model $A = (\mathcal{A}, \alpha)$ where \mathcal{A} does not only consist of \mathbb{N} , but also contains additional elements \square and \diamond . Then define $\alpha_{\text{plus}}(n, m)$ such that it models addition for $n, m \in \mathbb{N}$, but behaves differently if n or m are from $\{\square, \diamond\}$.

Exercise 2 (Matching):
(2 + 2 = 4 points)

- a) Consider the following pairs of terms s and t over the signature $\Sigma = \Sigma_0 \cup \Sigma_2$ with $\Sigma_0 = \{a\}$ and $\Sigma_2 = \{f\}$. Moreover, we have $\{x, y, z\} \subset \mathcal{V}$ for the set of variables \mathcal{V} . If s matches t , then give a suitable matcher σ . Otherwise give a brief (at most two sentences) explanation why there is no matcher.

1. $s = f(y, z), t = f(a, x)$
2. $s = f(x, a), t = f(a, x)$
3. $s = f(y, y), t = f(a, x)$
4. $s = f(x, y), t = f(f(x, y), a)$

- b) Let \sim be the matching relation, i.e., for two terms s and t we have $s \sim t$ iff s matches t .

Prove or disprove the following propositions:

1. For all terms s and t we have $s \sim t$ iff $t \sim s$ (i.e., the matching relation is symmetric).
2. For all terms s and t we have $s \sim t$ and $t \sim s$ iff $s = t$.

Exercise 3 (Stability):

(1 + 1 + 1 + 1 = 4 points)

Consider the following relations $\sim_1, \dots, \sim_4 \subseteq \mathcal{T}(\Sigma, \mathcal{V}) \times \mathcal{T}(\Sigma, \mathcal{V})$. Prove or disprove for each of these relations that they are stable.

- a) $s \sim_1 t$ iff the numbers of different variables in s and t are equal, i.e., $|\mathcal{V}(s)| = |\mathcal{V}(t)|$.
- b) $s \sim_2 t$ iff s is a subterm of t
- c) $s \sim_3 t$ iff s matches t
- d) $s \sim_4 t$ iff $\mathcal{V}(s) \subseteq \mathcal{V}(t)$

Hints:

- You can use the lemma proven in Exercise 4.

Exercise 4 (Induction):

(4 points)

Let $t \in \mathcal{T}(\Sigma, \mathcal{V})$, $\pi \in \text{Occ}(t)$ and $\sigma \in \text{SUB}(\Sigma, \mathcal{V})$. Show by induction over π that $(t|_{\pi})\sigma = (t\sigma)|_{\pi}$ holds.

Hints:

- In the induction base, prove the proposition for $\pi = \epsilon$.
- In the induction step, consider the case $\pi = i\pi'$, where as induction hypothesis, you can assume that $(q|_{\pi'})\mu = (q\mu)|_{\pi'}$ for all $q \in \mathcal{T}(\Sigma, \mathcal{V})$ and all $\mu \in \text{SUB}(\Sigma, \mathcal{V})$.